

# Functionally Constrained Algorithm Solves Convex Simple Bilevel Problems

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 $\mathcal{X}_{q}^{*}$ : minimizers of g(x)

 $f^*$ : minimal of f(x) over  $\mathcal{X}_a^*$ 

 $g^*$ : minimal of g(x)

• Our problem setup [Simple Bilevel Optimization]:

 $\min_{\mathbf{x}\in\mathcal{Z}} f(\mathbf{x}) \text{ s.t. } \mathbf{x} \in \mathcal{X}_g^* = \arg\min_{\mathbf{z}\in\mathcal{Z}} g(\mathbf{z}) \quad (1)$ 

- Minimize the upper-level objective over the solution set of a lower-level problem.
- $\boldsymbol{z}$ : feasible set; convex & compact with diameter D.
- f and g: upper-level and lower-level objective functions.
  - Assumption 1: f, g are convex and  $L_f$ ,  $L_g$ -smooth functions.
  - Assumption 2: f, g are convex and  $C_f$ ,  $C_g$ -Lipschitz continuous functions.
- - Lifelong learning, lexicographic optimization...
- Challenge:  $\mathcal{X}_g^*$  is not explicitly given.
  - Thus methods for constrained problems projected gradient method and Frank-Wolfe method are not applicable.





- **Our contribution:** •
  - Fundamental Difficulty of Simple BiO problems: Prove the intractability of any zero-٠ respecting first-order methods to find absolute optimal solutions.
  - **Near-Optimal Methods**: Propose a novel method with near-optimal rates for finding ٠ weak optimal solutions in both nonsmooth and smooth Simple BiO problems.

- Absolute optimal solution:  $|f(\hat{x}) f^*| \le \epsilon_f, g(\hat{x}) g^* \le \epsilon_g$ . Weak optimal solution:  $f(\hat{x}) f^* \le \epsilon_f, g(\hat{x}) g^* \le \epsilon_g$ .

### **I**. Hardness result: absolute optimal solution is not obtainable **M**

• Our result: It is generally **intractable** for any *zero-respecting first-order method* to absolute optimal solutions.

**Theorem 4.1.** For any first-order algorithm  $\mathcal{A}$  satisfying Assumption 3.4 that runs for T iterations and any initial point  $\mathbf{x}_0$ , there exists a (1, 1)-smooth instance of Problem (1) such that the optimal solution  $\mathbf{x}^*$  satisfies  $\|\mathbf{x}_0 - \mathbf{x}^*\|_2 \leq 1$  and  $\|f(\mathbf{x}_0) - f^*\| \geq \frac{1}{48}$ . For the iterates  $\{\mathbf{x}_k\}_{k=0}^T$  generated by  $\mathcal{A}$ , the following holds:

 $f(\mathbf{x}_k) = f(\mathbf{x}_0), \quad \forall 1 \le k \le T.$ 

**Theorem 4.2.** For any first-order algorithm  $\mathcal{A}$  satisfying Assumption 3.4 that runs for T iterations and any initial point  $\mathbf{x}_0$ , there exists a (1,1)-Lipschitz instance of Problem (1) and some adversarial subgradients  $\{\partial f(\mathbf{x}_k), \partial g(\mathbf{x}_k)\}_{k=0}^{T-1}$  such that the optimal solution  $\mathbf{x}^*$  satisfies  $\|\mathbf{x}_0 - \mathbf{x}^*\|_2 \leq 1$  and  $|f(\mathbf{x}_0) - f^*| \geq \frac{1}{4}$ . For the iterates  $\{\mathbf{x}_k\}_{k=0}^{T}$  generated by  $\mathcal{A}$ , the following holds  $f(\mathbf{x}_k) = f(\mathbf{x}_0), \quad \forall 1 < k < T.$ 

zero-respecting first-order method :  $\mathcal{A}$  generates test points  $\{\mathbf{x}_t\}_{t\geq 0}$  with  $\operatorname{supp}(\mathbf{x}_{t+1}) \subseteq \operatorname{supp}(\mathbf{x}_0) \cup \left(\bigcup_{0\leq s\leq t} \operatorname{supp}(\partial f(\mathbf{x}_s)) \cup \operatorname{supp}(\partial g(\mathbf{x}_s)))\right)$ 

## **II.** Hardness result: absolute optimal solution is not obtainable 道道大学

- Proof idea: we need to construct a "hard case".
  - Key concept: "first-order zero-chain" (Definition 3.1)
    - Applying zero-respecting first-order method to a first-order zero-chain with zero initialization: only one component of  $x_k$  becomes non-zero in each iteration.

#### . Near-optimal method for finding weak-optimal solutions

- Due to the intractability of obtaining absolute optimal solutions, we focus on proposing first-order methods for finding weak-optimal solutions: *f*(*x̂*) − *f*\* ≤ *ϵ<sub>f</sub>*, *g*(*x̂*) − *g*\* ≤ *ϵ<sub>g</sub>*.
- Step1: reformulate the original simple BiO problem to a functionally constrained problem.

$$\min f(x), s.t. \tilde{g}(x) \coloneqq g(x) - \hat{g}^* \le 0 \quad (2)$$

where  $\hat{g}^*$  is an approximation of the lower-level problem's optimal value  $g^*$ 

• **Step2**: Reduce Problem (2) to finding the smallest root of an auxiliary function (3), whose function value is defined by a discrete minimax problem. Such reformulation is introduced in Nesterov's *Lectures on convex optimization*.

$$\psi^*(t) \coloneqq \min_{x \in \mathcal{Z}} \{ \psi(t, x) \coloneqq \max\{ f(x) - t, \tilde{g}(x) \} \}$$
(3)

#### **II**. Near-optimal method for finding weak-optimal solutions



• To solve the smallest root of  $\psi^*(t)$ , we adopt a **bisection procedure**, and uses a **first-order subroutine**  $\mathcal{M}$  to estimate the function value of  $\psi^*(t)$  for a given t.

Algorithm: Functionally Constrained Bilevel Optimizer (FC-BiO) **Require:** desired accuracy  $\epsilon$ , total number of iterations T, initial bounds  $\ell, u$ , and first-order subroutine  $\mathcal{M}$ . Set  $N = \left[ \log_2 \frac{u-\ell}{\epsilon/2} \right]$ , K = T/N. Set  $\bar{\mathbf{x}} = \mathbf{x}_0$ . for  $k = 0, \cdots, N - 1$  do Set  $t = \frac{\ell + u}{2}$ . Solve with the subroutine  $(\hat{\mathbf{x}}_{(t)}, \hat{\psi}^*(t)) = \mathcal{M}(\bar{\mathbf{x}}, t, K)$ . Set  $\bar{\mathbf{x}} = \hat{\mathbf{x}}_{(t)}$ . if  $\hat{\psi}^*(t) \geq \frac{\epsilon}{2}$  then set  $\ell = t$ . else set u = t. End for **Return**  $\hat{\mathbf{x}} = \hat{\mathbf{x}}_{(u)}$  as the approximate solution.

. Near-optimal method for finding weak-optimal solutions

**first-order subroutine**  $\mathcal{M}$ :  $\min_{x \in \mathbb{Z}} \{ \psi(t, x) \coloneqq \max\{f(x) - t, \tilde{g}(x)\} \}$ 

Lipschitz objectives:

**Subgradient Method** 

$$\mathbf{x}_{k+1} = \Pi_{\mathcal{Z}} \big( \mathbf{x}_k - \eta \partial_{\mathbf{x}} \psi(t, \mathbf{x}_k) \big)$$

smooth objectives:

**Generalized Accelerated Gradient Method** 

$$\begin{aligned} \mathbf{x}_{k+1} &= \arg\min_{x\in\mathcal{Z}} \max \left\{ f(y_k) + \langle \nabla f(y_k), x - y_k \rangle + \frac{L}{2} \|x - y_k\|_2^2 - t \\ \tilde{g}(y_k) + \langle \nabla \tilde{g}(y_k), x - y_k \rangle + \frac{L}{2} \|x - y_k\|_2^2 \right\} \end{aligned}$$

 $x_{k+1}$  can be further written in the form of a projection. (Proposition 5.2)

#### **III**. Near-optimal method for finding weak-optimal solutions

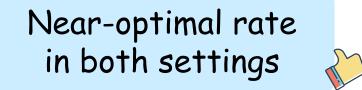


• Convergence rate of our FC-BiO method (Theorem 5.3, 5.4):

• Lipschitz case:  

$$\tilde{O}\left(\max\left\{\frac{C_f^2}{\epsilon_f^2}, \frac{C_g^2}{\epsilon_g^2}\right\}D^2\right)$$
• Smooth case:  

$$\tilde{O}\left(\max\left\{\sqrt{\frac{L_f}{\epsilon_f}}, \sqrt{\frac{L_g}{\epsilon_g}}\right\}D\right)$$
Where *D* is the diameter of *Z C\_f*, *C\_g*, *L\_f*, *L\_g* are Lipschitz/smooth constants, and  $\tilde{O}$  hides logarithmic terms







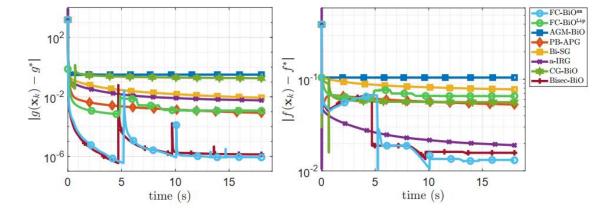


Figure 1: The performance of Algorithm 1 compared with other methods in Problem (11).

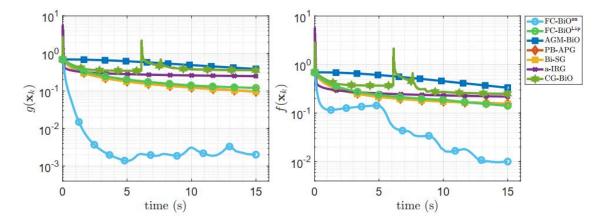


Figure 2: The performance of Algorithm 1 compared with other methods in Problem (12)

$$f(\mathbf{x}) = \frac{1}{2} \|\mathbf{x}\|_2^2, \quad g(\mathbf{x}) = \frac{1}{2} \|A\mathbf{x} - b\|_2^2,$$

minimum norm solution of Linear Regression. 400 datapoints. 700+ features.

$$f(\mathbf{x}) = \frac{1}{m} \sum_{i=1}^{m} \log(1 + \exp(-(A_i^{val})^{\top} \mathbf{x} \mathbf{b}_i^{val})),$$
$$g(\mathbf{x}) = \frac{1}{m} \sum_{i=1}^{m} \log(1 + \exp(-(A_i^{tr})^{\top} \mathbf{x} \mathbf{b}_i^{tr})).$$

Overparameterized Logistic Regression 10000 datapoints, 40000+ features.



