



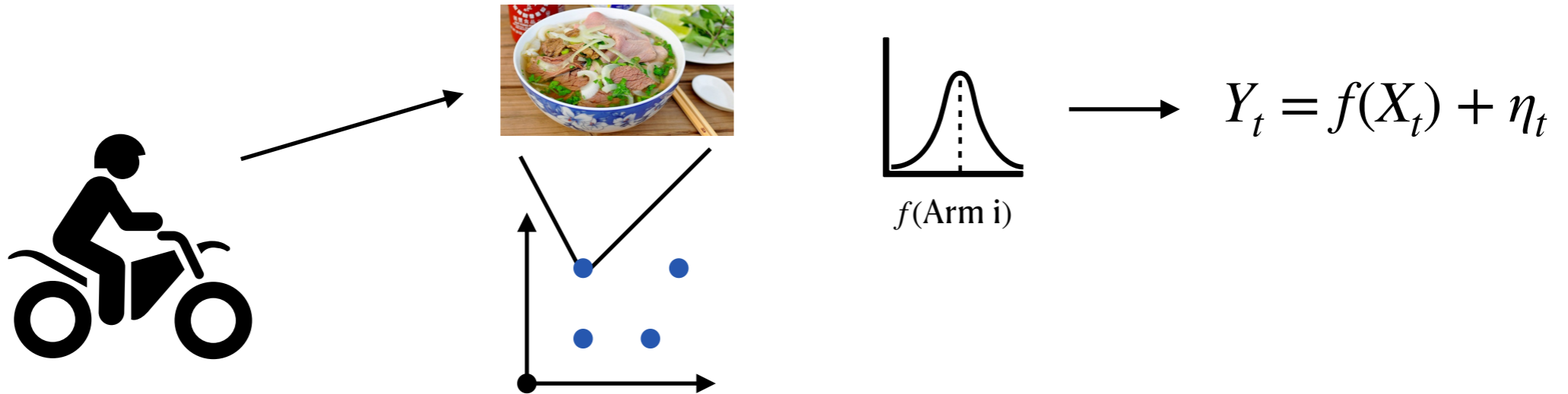
Symmetric Linear Bandits with Hidden Symmetry

Nam Phuong Tran¹, The-Anh Ta²,
Debmalya Mandal¹, Long Tran-Thanh¹

University of Warwick¹, Data 61 Australia²



Linear Bandits



Linear Bandit: $y_t = f(x_t) + \eta_t$, $f(x_t) = \langle x_t, \theta \rangle$. $\theta \in \mathbb{R}^d$, $\eta \sim N(0, \sigma)$.

The goal is to minimise the expected regret $\mathbf{R}_T = \mathbb{E} \left[\sum_{t=1}^T f_{\star} - f(X_t) \right]$.

Minimax regret: $\mathbf{R}_T = \Theta(d\sqrt{T})$

Linear Bandits and The Curse of Dimensionality

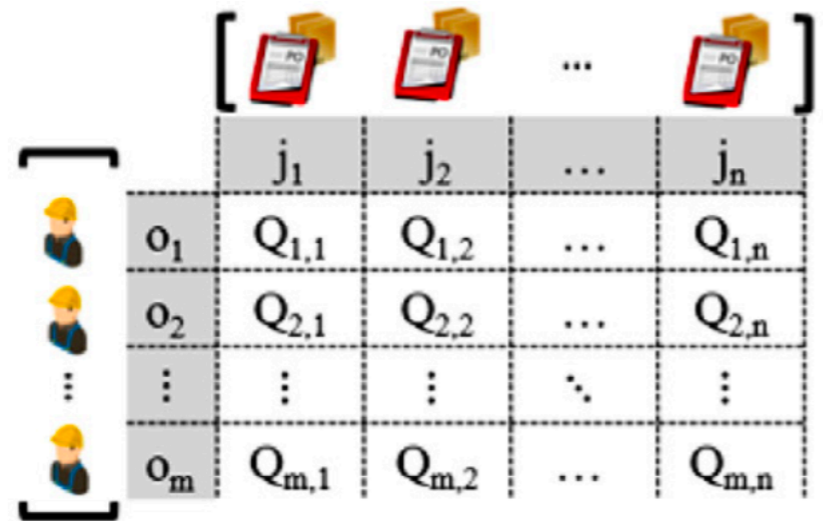
Linear Bandit: $y_t = f(x_t) + \eta_t$, $f(x_t) = \langle x_t, \theta \rangle$. $\theta \in \mathbb{R}^d$, $\eta \sim N(0, \sigma)$.







- Minimax regret: $\mathbf{R}_T = \Theta(d\sqrt{T})$
- Curse of dimensionality: d is large, vacuous regret.

Inductive bias: θ lies in low-dimensional structures.

- **Sparsity** $\|\theta\|_0 \leq s$: $\mathbf{R}_T = O(s \log(d)\sqrt{T})$
- **Question:** Is there any other low-dimensional structure that help overcome the curse of dimensionality?

Hint from supervised learning



				...	
		j_1	j_2	...	j_n
	o_1	$Q_{1,1}$	$Q_{1,2}$...	$Q_{1,n}$
	o_2	$Q_{2,1}$	$Q_{2,2}$...	$Q_{2,n}$
⋮	⋮	⋮	⋮	⋮	⋮
	o_m	$Q_{m,1}$	$Q_{m,2}$...	$Q_{m,n}$



What is symmetry? $f(g \cdot x) = f(x)$.

invariance

Break the curse of dimensionality with *partial* knowledge of the symmetry structure?

Optimal regret bound: What is the optimal regret bound with the presence of symmetry.

Symmetric Linear Bandit

Linear Bandit: $y_t = f(x_t) + \eta_t$, $f(x_t) = \langle x_t, \theta \rangle$.

- $x \in \mathcal{X} \subset \mathbb{R}^d$; $\theta \in \mathbb{R}^d$, $\eta \sim N(0, \sigma)$, where $d \geq T$.
- **Symmetry:** $\mathcal{G} \leq \mathcal{S}_d$ be a subgroup of symmetric group of set $[d] = \{1, \dots, d\}$.
- **Group action:** For $g \in \mathcal{G}$, let $g \cdot x$ permute the coordinate of x .
- **Invariant mean reward:** $\forall x \in \mathcal{X}, g \in \mathcal{G}; \langle x, \theta \rangle = \langle g \cdot x, \theta \rangle$.
- **Hidden symmetry:** \mathcal{G} is only known to be a subgroup of \mathcal{S}_d .

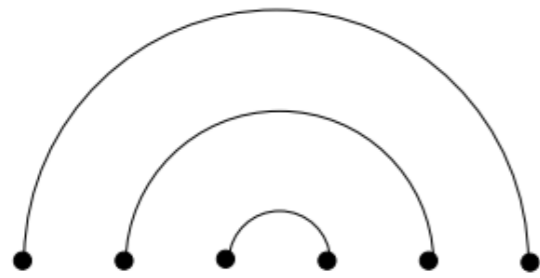
Impossibility Result

Proposition 1 (Impossibility result): Assume, \mathcal{X} is unit cube, and f is invariant w.r.t. action of \mathcal{G} , s.t., $\dim(\text{Fix}_{\mathcal{G}}) = 2$, any bandit algorithm suffer regret $\Omega(d\sqrt{T})$.

Extra assumptions needed!!!

Assumption 2 (Sub-exponential number of partitions). The partition corresponding to \mathcal{G} belongs to a small subclass of partitions $\mathcal{Q}_{d, \leq d_0} \subset \mathcal{P}_{d, \leq d_0}$. Moreover, for any $k \leq d_0$, there exist a constant $c > 0$ such that $|\mathcal{Q}_{d, k}| \leq O(d^{cd_0})$.

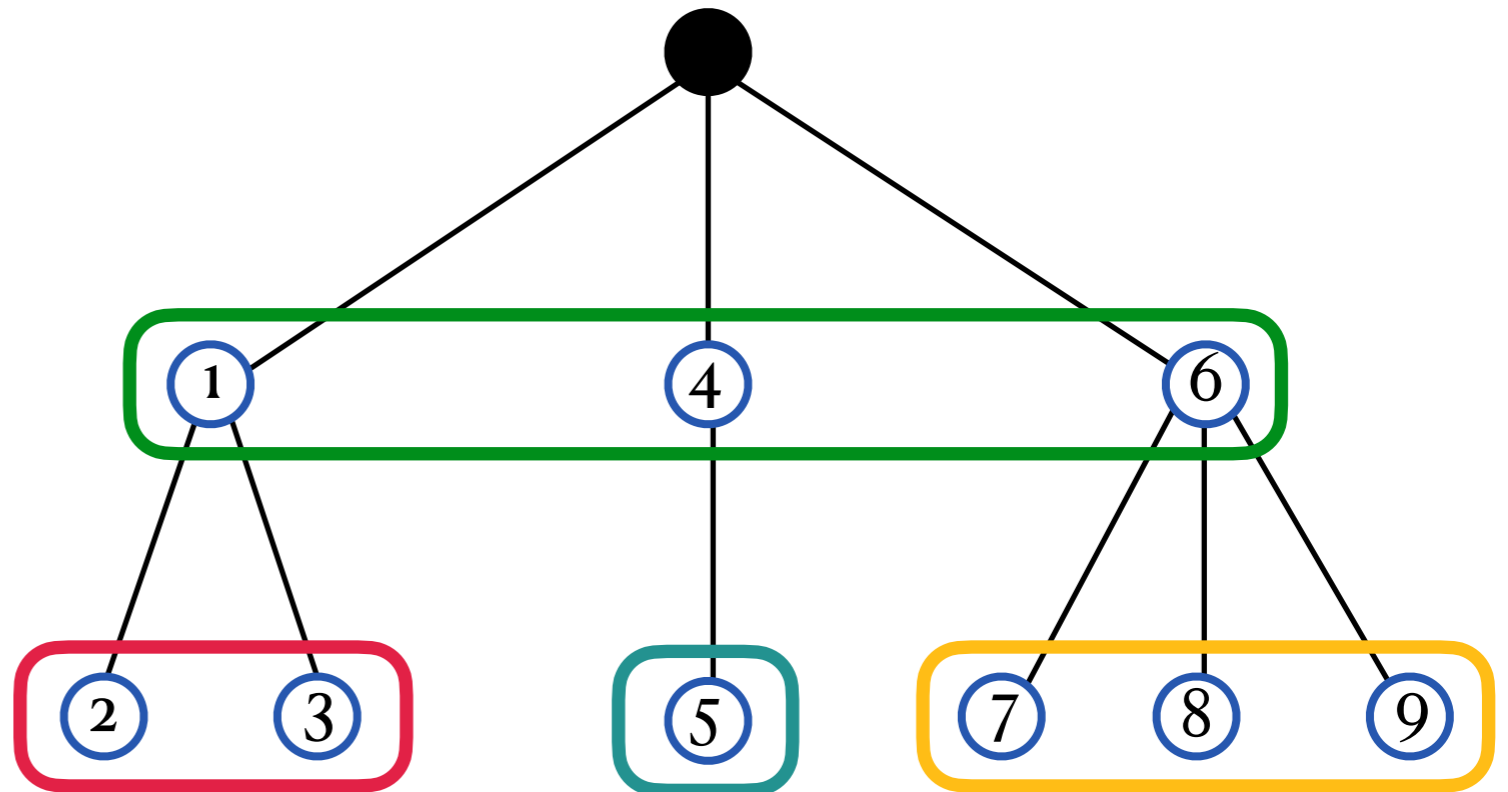
Sub-exponential collection of models



Non-crossing Partitions



Non-nesting Partitions



Partitions respect hierarchy

Assumption 2 (Sub-exponential number of partitions). The partition corresponding to \mathcal{G} belongs to a small subclass of partitions $\mathcal{Q}_{d, \leq d_0} \subset \mathcal{P}_{d, \leq d_0}$. Moreover, for any $k \leq d_0$, there exist a constant $c > 0$ such that $|\mathcal{Q}_{d, k}| \leq O(d^{cd_0})$.

A Lower Bound of Symmetric Linear Bandits

Recall the equivalence between interval partition and sparsity!

Regret Lower Bound: There exist symmetric linear bandit instances in which Assumption 1 holds, such that any bandit algorithm must suffer regret

$$\mathbf{R}_T = \Omega \left(\min \left(C_{\min}^{-1/3}(\mathcal{X}) d_0^{2/3} T^{2/3}, \sqrt{dT} \right) \right)$$

Regret Upper Bound

Theorem 7 (Regret upper bound): with the choice of

$t_1 \cong C_{\min}(\mathcal{X})^{-1/3} K_x^{1/3} d_0^{1/3} T^{2/3}$, the regret of Algorithm 1 is upper bounded as

$$\mathbf{R}_T = O\left(C_{\min}^{-1/3}(\mathcal{X}) d_0^{2/3} T^{2/3} \log(dT)^{1/3}\right)$$

Theorem 1 (Regret upper bound): With well-separateness assumption, Algorithm returns $\widehat{m} \ni \theta_\star$ with probability at least $1 - 1/T$, and its regret is upper bounded as

$$\mathbf{R}_T = O\left(d_0 \sqrt{T \log(d_0 T)}\right)$$

Conclusion

- We show impossibility result $\Omega(d\sqrt{T})$ for general symmetry.
- With cardinality constraint, it possible to achieve regret $\tilde{\Theta}(d_0^{2/3}T^{2/3})$, and $\tilde{\Theta}(d_0\sqrt{T})$.