# Rényi-infinity Uniform Sampling via Algorithmic Diffusion

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# Outline

• Part I: Uniform sampling over a convex body through diffusions - arXiv:2405.01425

• Part II: Warm-start generation without "TV collapse"

- arXiv:2407.12967



# Part I - Collaborators

## In-and-Out: Algorithmic Diffusions for Sampling Convex Bodies NeurIPS'24



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# Uniform sampling is (maybe) all you need

Sampling from a log-concave  $d\pi \propto \exp(-V) dx$ 



- Main subroutine in volume computation
- System biology

. . .

Uniform sampling from a convex body *K* 



# Uniform sampling in formulation

**Problem**. Let *K*: convex body in  $\mathbb{R}^d$  and  $\pi = \text{Unif}(K)$ . How many <u>membership oracle</u> queries are needed to generate a sample *X* whose law is  $\varepsilon$ -close to  $\pi$  in some *D*?

 $D(law(X), \pi) \leq \varepsilon$  for some probability divergence/distance D = TV, KL,  $\chi^2$  etc.



# Geometric random walk

Ball walk $(\delta)$ 

- 1. Pick  $z \in B_{\delta}(x)$
- 2. Move to z if  $z \in K$ . Stay at x o.w.

## Hit-and-Run

- 1. Pick a uniform random line  $\ell$  through the current point x
- 2. Move to a uniform random point on the chord  $\ell \cap K$



# Another line of log-concave sampling research

## **Problem (Well-conditioned log-concave sampling)**.

(strong convexity and smoothness of a potential V) over  $\mathbb{R}^d$ .

whose law is  $\varepsilon$ -close to  $\pi$ ?

- Let  $\pi \propto \exp(-V)$  be a smooth unconstrained distribution with  $\alpha I \leq \nabla^2 V \leq \beta I$
- How many access to the first-order oracle of V are needed to generate a sample X



# Well-conditioned log-concave sampling

- General approach for getting an implementable algorithms 1. Understand the Langevin dynamics (SDE) with stationary  $\pi \propto \exp(-V)$ :  $dX_t = -\nabla V(X_t) dt + \sqrt{2} dB_t$
- 2. Discretize it in time:
  - Euler-Maruyama discretization
  - Randomized midpoint method
  - So on...



# Well-conditioned log-concave sampling

## Analysis

- 1. Establish the mixing of the Langevin dynamics in  $W_2$ , KL,  $\chi^2$  (or generally Rényi)
- 2. Discretization-analysis somehow preserves the mixing metric
  - Girsanov's theorem [Dalalyan and Tsybakov'12]
  - Interpolation method [Vempala and Wibisono'19]
  - Hypercontractivity [Chewi et al.'21]
  - Shifted composition rule [Altschuler and Chewi'24]



# Hierarchy of probability distance/divergence



$$\mathbf{TV}(\mu, \pi) = \sup_{S} |\mu(S) - \pi(S)|$$

 $W^{2}(\mu, \pi) = \inf_{\Gamma(\mu, \pi)} \mathbb{E}_{(X, Y) \sim \Gamma}[\|X - Y\|^{2}]$ 

# A current state of affairs

## [Constrained sampling]

Algs Ball walk, Hit-and-Run Metrics TV,  $\chi^2$ Tools Conductance

## Fundamental gap here?

## [Unconstrained sampling] Algs Langevin-based Metrics $\mathscr{R}_{q}$ Tools

Wasserstein calculus, optimal transport, Markov semigroup theory, interpolation method, Girsanov's argument, Shifted composition rule,.....



# Let's bridge this gap

## [Constrained sampling]

Algs <u>New sampler</u> Metrics

$$\mathscr{R}_q$$
 (and  $\mathscr{R}_\infty$  in fact)

Tools

Continuous interpolation via a forward/backward SDE

Can borrow these techniques!

## [Unconstrained sampling]

## Algs Langevin-based Metrics

## $\mathscr{R}_q$ Tools

Wasserstein calculus, optimal transport, Markov semigroup theory, interpolation method, Girsanov's argument, Shifted composition rule,.....



# In-and-Out



In-and-Out

[Forward] Sample  $y_{i+1} \sim N(x_i, hI_d)$ [Backward] Sample  $x_{i+1} \sim N(y_{i+1}, hI_d)|_K$ 

\* One iteration = forward + backward step



# In-and-Out

**Input:** initial point  $x_0 \sim \pi_0$  & convex body  $K \subset \mathbb{R}^d$  & threshold N & step size h **Output:**  $X_T$ 

• For 
$$i = 0, ..., T$$

- 1. Sample  $y_{i+1} \sim N(x_i, hI_d) = x_i + N(0, hI_d)$
- 2. Sample  $x_{i+1} \sim N(y_{i+1}, hI_d)|_{\kappa}$

[Implementation]

 $-x_{i+1} \sim N(y_{i+1}, hI_d)$  until  $x_{i+1} \in K$ 

- If [# attempts  $\geq N$ ], then declare **Failure** 

Where does it come from? Connection to proximal sampler [Lee, Shen, and Tian'21] **Goal:** Sample from  $\pi(x) \propto \exp(-V(x))$  over  $\mathbb{R}^d$ To this end, augment another variable  $y \in \mathbb{R}^d$  to consider  $\pi(x, y) \propto \exp(-$ 

**Algorithm**: Repeat

1. Sample  $y_{i+1} \sim \pi^{Y|X=x_i}(y) = N(x_i, hI_d)$ 

2. Sample  $x_{i+1} \sim \pi^{X|Y=y_{i+1}}(x) \propto \exp(-V)$ 

$$-V(x) - \frac{1}{2h} ||x - y||^2$$

$$V(x) - \frac{1}{2h} \|x - y\|^2$$



# Where does it come from?

## Connection to proximal sampler [Lee, Shen, and Tian'21]

## **Algorithm**: Repeat

- 1. Sample  $y_{i+1} \sim \pi^{Y|X=x_i}(y) = N(x_i, hI_d)$
- 2. Sample  $x_{i+1} \sim \pi^{X|Y=y_{i+1}}(x) \propto \exp(-\frac{1}{2})$

## **In-and-Out** is the Proximal sampler with $\pi(x) \propto 1_K(x)$

$$\frac{1}{2h} \|x - y\|^2 \big|_K$$



# Proximal sampler in measure level





# Outline of analysis

- **1. Contraction** through one-iteration of INO (proximal sampler)
- 2. Query complexity of the implementation for the backward step

Forward / backward SDE interpretation by [Chen et al.'22]

Forward step: Sample  $y_{i+1} \sim \pi^{Y|X=x_i}(y) = N(x_i, hI_d)$ 

$$Z_0 \sim \text{law}(X_i)$$
  $dZ_t = dZ_t$ 





Forward / backward SDE interpretation by [Chen et al.'22]

Backward step: Sample  $x_{i+1} \sim \pi^{X|Y=y_{i+1}}(x)$ 

$$Z_0 \sim \text{law}(X_i) \qquad \text{d}Z_t = \text{d}B_t \text{ for } t \in [0,h] \qquad Z_h \sim \text{law}(Y_{i+1})$$

$$dZ_t^{\leftarrow} = \nabla \log(\pi^X Q_{h-t})(Z_t^{\leftarrow}) \, \text{d}t + \text{d}B_t \text{ for } t \in [0,h] \qquad Z_0^{\leftarrow} = Z_h$$

$$Z_0 \sim \text{law}(X_i) \qquad \text{d}Z_t = \text{d}B_t \text{ for } t \in [0,h] \qquad Z_h \sim \text{law}(Y_{i+1})$$

$$Z_h^{\leftarrow} \sim \text{law}(X_{i+1}) \qquad \text{d}Z_t^{\leftarrow} = \nabla \log(\pi^X Q_{h-t})(Z_t^{\leftarrow}) \text{ d}t + \text{d}B_t \text{ for } t \in [0,h] \qquad Z_0^{\leftarrow} = Z_h$$

$$x) \propto \exp\left(-V(x) - \frac{1}{2h} \|x - y_{i+1}\|^2\right)$$



**A1.** Avoid discretization-analysis A2. Use tools from Markov semigroup theory

**Q.** Benefits of introducing an SDE representation?

**Contraction via forward/backward heat-flow** Forward / backward SDE interpretation by [Chen et al.'22] Using Markov semigroup theory + Wasserstein calculus, one can show Forward SDE:  $dZ_t = dB_t$  with  $Z_0 \sim \mu_i^X \implies Z_h \sim \mu_{i+1}^Y$ 

$$\chi^{2}(\mu_{i+1}^{Y} \| \pi^{Y}) \leq \frac{\chi^{2}(\mu_{i}^{X} \| \pi^{X})}{1 + h/C_{\mathsf{PI}}(\pi^{X})}$$
$$\mathsf{KL}(\mu_{i+1}^{Y} \| \pi^{Y}) \leq \frac{\mathsf{KL}(\mu_{i}^{X} \| \pi^{X})}{1 + h/C_{\mathsf{LSI}}(\pi^{X})}$$

 $\frac{1}{X_{Y}}$  under a Poincaré inequality (PI)

- under a log-Sobolev inequality (LSI)

**Contraction via forward/backward heat-flow** Forward / backward SDE interpretation by [Chen et al.'22] Using Markov semigroup theory + Wasserstein calculus, one can show Backward SDE:  $dZ_t^{\leftarrow} = \nabla \log(\pi^X Q_{h-t})(Z_t^{\leftarrow})$ 

$$\chi^{2}(\mu_{i+1}^{X} \| \pi^{X}) \leq \frac{\chi^{2}(\mu_{i+1}^{Y} \| \pi^{Y})}{1 + h/C_{\mathsf{PI}}(\pi^{X})}$$
$$\langle \mathsf{L}(\mu_{i+1}^{X} \| \pi^{X}) \leq \frac{\mathsf{KL}(\mu_{i+1}^{Y} \| \pi^{Y})}{1 + h/C_{\mathsf{LSI}}(\pi^{X})}$$

$$(t^{\leftarrow}) dt + dB_t \text{ with } Z_0^{\leftarrow} \sim \mu_{i+1}^Y \Longrightarrow Z_h^{\leftarrow} \sim \mu_{i+1}^X$$

 $\frac{1}{X_{1}}$  under a Poincaré inequality (PI)

under a log-Sobolev inequality (LSI)



## Forward / backward SDE interpretation by [Chen et al.'22]

Composing one forward + backward contraction,

$$\chi^{2}(\mu_{i+1}^{X} \| \pi^{X}) \leq \frac{\chi^{2}(\mu_{i}^{X} \| \pi^{X})}{\left(1 + h/C_{\mathsf{PI}}(\pi^{X} \| \pi^{X}) + KL(\mu_{i+1}^{X} \| \pi^{X})\right)}$$

$$\mathsf{KL}(\mu_{i+1}^{X} \| \pi^{X}) \leq \frac{\mathsf{KL}(\mu_{i}^{X} \| \pi^{X})}{\left(1 + h/C_{\mathsf{LSI}}(\pi^{X})\right)}$$

\* Can still use this result, though  $\pi_K$  is not smooth around  $\partial K$ ! (Convolve with  $N(0, \epsilon I_d)$  & Use the lower-semi continuity of f-divergence as  $\epsilon \to 0$ )

\* Can be extended to a q-Rényi divergence

 $\frac{1}{(N)^2}$  under a Poincaré inequality (PI)

 $\frac{1}{\sqrt{2}}$  under a log-Sobolev inequality (LSI)

Functional inequalities for  $\pi$ 

Poincaré inequality (PI)

Log-Sobolev inequality (LSI)

**Caveat**. This depends only on a measure  $\pi$ , not on a Markov chain/kernel

 $\operatorname{var}_{\pi} f \leq C_{\operatorname{PT}}(\pi) \mathbb{E}_{\pi}[\|\nabla f\|^2]$  for any smooth  $f: \mathbb{R}^d \to \mathbb{R}$ 

 $\operatorname{Ent}_{\pi}(f^2) \leq 2C_{\operatorname{LST}}(\pi) \mathbb{E}_{\pi}[\|\nabla f\|^2]$  for any smooth  $f: \mathbb{R}^d \to \mathbb{R}$ 



Known results on log-concave distribution  $\pi \propto \exp(-V)$  over  $\mathbb{R}^d$ 

- 1.  $C_{\text{PT}}(\pi) \leq C_{\text{IST}}(\pi)$  in general
- 2.  $C_{\rm PT}(\pi)$ 
  - $\|\operatorname{Cov}(\pi)\|_{\operatorname{OD}} \leq C_{\operatorname{PI}}(\pi) \leq \psi_d \cdot \|\operatorname{Cov}(\pi)\|_{\operatorname{OD}}$
  - KLS conjecture:  $\psi_d = \Theta(1)$
  - $\psi_d \leq \log d$  [Klartag'23]

3.  $C_{LSI}(\pi) = O(D^2)$  for a log-concave  $\pi$  with support of diameter D





Relating mixing guarantees to functional inequalities

**Theorem.** For  $\varepsilon \in (0,1)$  and  $K \subset B_D(0)$ , INO with step-size h and M-warm initial distribution achieves  $\mathscr{R}_q(\mu_n || \pi_K) \leq \varepsilon$  after the following # of iterations:

$$n \asymp \min \begin{cases} qd^2 \|\operatorname{Cov}(\pi_K)\|_{\operatorname{Op}} \log \frac{M}{\varepsilon} & \text{for } q \ge 2\\ qd^2 D^2 \log \frac{\log M}{\varepsilon} & \text{for } q \ge 1 \end{cases}$$

(  $\uparrow$  ) Substitute the known bounds on  $C_{\rm PT}, C_{\rm LST}$  and  $h \asymp d^{-2}$ 



Rejection sampling for the backward step  $x_{i+1} \sim N(y_{i+1}, hI_d)|_{K}$ 

[Implementation via rejection sampling]

 $-x_{i+1} \sim N(y_{i+1}, hI_d)$  until  $x_{i+1} \in K$ 



## But this is bound to fail



## Rejection sampling for the backward step $x_{i+1} \sim N(y_{i+1}, hI_d)|_{K}$

Suppose we're already at stationarity  $\pi^X = \text{Unif}(K)$ 

 $\rightarrow \pi^Y = \pi^X$ 



 $\rightarrow$  Simply, the success probability of the rejection sampling at y

\* N(0,
$$hI_d$$
) =  $\frac{\ell(y)}{\operatorname{vol}(K)}$ 

where  $\ell(y)$  is a Gaussian version of local conductance [Kannan et al.'97] defined by

$$\frac{\exp(-\frac{1}{2h}\|x-y\|^2)\,\mathrm{d}y}{\exp(-\frac{1}{2h}\|x-y\|^2)\,\mathrm{d}y}$$



## Rejection sampling for the backward

Then the expected number of trials (until success) for one iteration is

$$\mathbb{E}_{\pi^{Y}}\left[\frac{1}{\ell(y)}\right] = \int_{\mathbb{R}^{d}} \frac{1}{\ell(y)} \frac{\ell(y)}{\operatorname{vol}(K)} \, \mathrm{d}y = \infty$$

**Q**. Can bypass this issue?

$$\mathsf{step} \ x_{i+1} \sim \mathsf{N}(y_{i+1}, hI_d) \big|_K$$



## **Rejection sampling for the backward step** $x_{i+1} \sim N(y_{i+1}, hI_d)|_{K}$



**Lemma**.  $\pi^{Y}(\mathbb{R}^{d}\setminus K_{\delta}) \leq \exp(-\Theta(t^{2}))$  for step-size  $h = \Theta(d^{-2})$  and  $\delta = t/d$ .

 $\rightarrow K_{\delta}$  is sort of "effective domain" of  $\pi^{Y}$ 





**Insight**: Ignore whatever happens outside of this effective domain  $K_{\delta}$ 

## **Rejection sampling for the backward**

**Proposition**. For  $y \in K_{\delta}^{c}$  with  $\delta = t/d$  and  $h \asymp d^{-2}$ ,

- Expected #trials from  $K_{\delta}^c$  for the rejection sampling  $\rightarrow \ell(y)^{-1} \ge \exp(\Omega(t^2))$
- $\therefore$  Can ignore algorithmic behaviors from  $K^c_{\delta}$  by setting a threshold  $N = O(\exp(t^2))$ and considering the algorithm as having "failed" if #trials  $\geq N$

$$\mathsf{step} \ x_{i+1} \sim \mathsf{N}(y_{i+1}, hI_d) \big|_K$$

- **Q.** Characteristic of the complement of the effective domain  $K_{\delta}$ ?

  - $\ell(y) \leq \exp(-\Omega(t^2))$ .



- **Rejection sampling for the backward**
- - 1. failure prob. of one backward-step  $\leq \delta/T$
  - 2. the expected # of queries per backward-step  $\leq M$  polylog $(TM/\delta)$

step 
$$x_{i+1} \sim \mathsf{N}(y_{i+1}, hI_d) |_K$$

**Theorem**. (Complexity of backward step) For failure prob.  $\delta \in (0,1)$  and  $T \in \mathbb{N}$ , there exists suitable choices of parameters h, N such that

 $\therefore$  During T iterations, (1) the total failure prob. is  $\leq \delta$ , and (2) the total query complexity is  $\widetilde{O}(MT \operatorname{polylog}(1/\delta))$ 



# **Guarantee of In-and-Out**

## **Assumption:**

- 1. Access to a membership oracle for a convex body  $K \subset \mathbb{R}^d$  with unit ball in K.
- 2. An initial  $\pi_0$  is *M*-warm w.r.t. target  $\pi = \text{Unif}(K)$ - Precisely,  $d\pi_0/d\pi \leq M$  a.s.

- **Theorem.** Given failure prob.  $\delta \in (0,1)$ , target acc.  $\varepsilon \in (0,1)$ , and  $q \ge 1$ , there exists choices of parameters h, N such that with probability  $\geq 1 - \delta$ ,
  - INO started at  $\pi_0$  ensures  $\mathscr{R}_q(\operatorname{law}(X_n) \| \pi) \leq \varepsilon$ after  $n = \widetilde{O}(qd^2 \| \operatorname{Cov}(\pi) \|_{\operatorname{Op}} \operatorname{polylog}(M/\delta \varepsilon))$  iterations, using  $\widetilde{O}(qMd^2 \| \text{Cov}(\pi) \|_{OD}$  polylog $(1/\delta \varepsilon)$  membership queries in expectation.





# Matching results of Ball walk

exists choices of parameters h, N such that with probability  $\geq 1 - \delta$ ,

Previous best complexity via **Ball walk**:

- **Theorem.** Given failure prob.  $\delta \in (0,1)$ , target acc.  $\varepsilon \in (0,1)$ , and  $q \ge 1$ , there
  - INO started at  $\pi_0$  ensures  $\mathscr{R}_q(\operatorname{law}(X_n) \| \pi) \leq \varepsilon$ by using  $\widetilde{O}(qMd^2 \| \operatorname{Cov}(\pi) \|_{\operatorname{Op}} \operatorname{polylog}(1/\delta \varepsilon))$  membership queries in expectation.

- $\rightarrow$  Achieving  $\varepsilon$ -TV distance from *M*-warm start needs  $O(Md^2 \|Cov(\pi)\|_{OD} \operatorname{polylog}(1/\delta\varepsilon))$  queries.
- INO recovers the matching result under stronger performance metrics and principled approaches!



# **Chicken-and-egg problem**

## Awkward situation...

Observation

# INO needs $\widetilde{O}(qMd^2 \| \text{Cov}(\pi) \|_{\text{op}} \text{ polylog} \frac{1}{\delta c})$ queries where $M = \exp(\mathscr{R}_{\infty}(\pi_0 \| \pi))$

## Needs $\mathscr{R}_{\infty}$ -warmness to get $\mathscr{R}_{q}$ -result (denote INO: $\mathscr{R}_{\infty} \to \mathscr{R}_{q}$ ). Same issue with BW (SW + rejection): $\mathscr{R}_{\infty} \to TV$ .

**Q**. How to get a warm-start in  $\mathscr{R}_{\infty}$ ?



# Part II - Collaborators

## **Rényi-infinity constrained sampling with** $d^3$ **membership queries** SODA'25



Matthew Zhang University of Toronto



## Warm-start generation

**Problem.** Let  $K \subset \mathbb{R}^d$  be a "well-rounded" convex body (i.e.,  $\mathbb{E}_{\pi}[||X||^2] = O(d)$ ) containing a unit ball. Can we generate a warm start X such that

 $\mathscr{R}_{\sim}(\operatorname{law}(X) \parallel \pi) = O(1)?$ 

**Note)** There is a known method for making K well-rounded [Jia et al.'21]



A common approach is "annealing":

$$\mu_0 \to \mu_1 \to \cdots \to \mu$$

- $\mu_0$ : easy dist. (from which we can easily sample)
- Generate  $\mu_{i+1}$  starting from  $\mu_i$  (by some samplers)
- Here,  $\mu_i$  is a warm start for  $\mu_{i+1}$  (i.e.,  $\mu_i$  and  $\mu_{i+1}$  are already close)

## $\mu_i \to \mu_{i+1} \to \cdots \to \mu_k \to \pi$



- 1. Uniform annealing [Dyer-Frieze-Kannan'89 ~ Kannan-Lovász-Simonovits'97]  $\mu_i = K \cap (2^{i/d} B_1(0))$  $\mathscr{R}_{\infty}(\mu_{i} \| \mu_{i+1}) = O(1)$
- 2. Exponential annealing [Lovász and Vempala'06]  $\mu_i \propto \exp(-a_i^{\mathsf{T}} x) |_{\mathbf{K}}$
- 3. Gaussian annealing [Cousin and Vempala'18]
  - $\mu_i \propto \exp\left(-\frac{1}{2\sigma^2} \|x\|^2\right)|_K$

## Recall $\mathscr{R}_2 = \log(\chi^2 + 1)$

 $\Re_{2}(\mu_{i} \| \mu_{i+1}) = O(1)$ 



sampling the annealing distribution  $\mu_i$ .

However, HAR or SW has guarantees in  $\chi^2$  or TV.

Previous works rely on Hit-and-Run or [Speedy walk + rejection sampling] for



 $P_1$ : Markov kernel of the MCMC sampler

**MCMC Sampler** started at  $\mu_0$  will output  $X \sim \mu_0 P$  with  $TV(\mu_0 P_1, \mu_1) \leq \varepsilon$ In the next phase, an initial dist. is **in fact**  $\mu_0 P_1$ , not  $\mu_1$ . No triangle inequality coupling TV and  $\mathscr{R}_{\infty}$ 





However, the annealing algorithm still proceeds as if the starting distribution is  $\mu_1$ Previous works use a <u>coupling argument</u> for analysis, reducing everything to TV.



- - Previous approaches cannot avoid this "TV-collapse" issue

ends up collapsing to TV as well

• Due to inexact error from Markov chains, any guarantee is eventually <u>collapsed to TV</u>

• If INO uses a warm start generated by this annealing scheme, then its final guarantee





If a sampler has a  $\mathscr{R}_{\infty}$ -guarantee, then can relay  $\mathscr{R}_{\infty}$ -guarantees through the triangle inequality



# $\mathscr{R}_{\infty}$ is difficult

- Prior sampling  $\mathscr{R}_q$ -guarantees involve a complexity at least linear in q
  - Useless for  $\mathscr{R}_{\infty}$
- A Markov-semigroup approach used for  $\mathscr{R}_q$  doesn't go through for  $\mathscr{R}_{\infty}$

## In this work, we boost $TV \rightarrow \mathscr{R}_{\infty}$ without overhead via a log-Sobolev inequality (LSI)



# **Revisit the theory of Markov semigroups**

• Let  $P: \Omega \times \mathscr{F} \to [0,1]$  be a Markov kernel. Then,

 $\mu P(\cdot) :=$ 

Pf(x) :=

• Convergence rate is characterized by the **contractivity** of a Markov kernel:

 $||P||_{I^{p} \to I^{p}}$ 

where  $L_0^p := \{f : \mathbb{E}_{\pi}[|f|^p] < \infty, \mathbb{E}_{\pi}f = 0\}.$ 

$$\int_{\Omega} P(\cdot | x) \mu(dx)$$
$$\int_{\Omega} f(y) P(dy | x)$$

$$:= \sup_{0 \neq f \in L_0^p} \frac{\|Pf\|_{L^p}}{\|f\|_{L^p}}$$



# Revisit the theory of Markov semigroups

The most classical setting is the " $L^2(\pi) \rightarrow L^2(\pi)$  contraction"

- If  $\gamma := ||P||_{L^2 \to L^2}$ , then the so-called **spectral gap** of P is  $1 \gamma$

**Q**. What about contraction in  $L^{\infty} \to L^{\infty}$ ?  $(\text{recall } ||f||_{L^{\infty}} := \inf \{C : |f| \le C\} = \text{esssup } |f|)$ 



# Revisit the theory of Markov semigroups

**Q**. What about contraction in  $L^{\infty} \to L^{\infty}$ ?

where  $1_{\pi}$  is the operator defined by  $1_{\pi}(f) := \mathbb{E}_{\pi}f$ .

- **Theorem** [Rudolf'11]. Let P be a Markov kernel reversible w.r.t. stationary  $\pi$ . Then,
  - $||P^n 1_{\pi}||_{L^{\infty} \to L^{\infty}} = 2 \operatorname{esssup}_{Y} \operatorname{TV}(\delta_{X} P^n, \pi)$



# Convergence from any start implies $L^{\infty}$ -contraction

By substituting  $f = \frac{d\mu}{d\pi} - 1$ , one can deduce  $\mathscr{R}_{\infty}(\mu P^{n} \| \pi) \leq \left\| \frac{\mathsf{d}(\mu P^{n})}{\mathsf{d}\pi} - 1 \right\|_{L^{\infty}}$  $\log(1+x) \le x$ 

$$\leq \left\| \frac{\mathrm{d}\mu}{\mathrm{d}\pi} - 1 \right\|_{L^{\infty}} \cdot 2 \operatorname{esssup}_{x} \operatorname{TV}(\delta_{x}P^{n}, \pi)$$

## $\therefore$ Uniform TV-bound over any start $x \in \Omega \implies \mathscr{R}_{\infty}$ -bound

# Functional inequalities for boosting

**Q**. When can we bound sup  $TV(\delta_x P^n, \pi)$  without huge overhead?  $x \in \Omega$ 

 $\rightarrow$  (**PI**) ensures an exponential contraction  $\chi^2(\delta_{\mathbf{x}}P^n \parallel \pi) \lesssim \exp$ 

 $\rightarrow$  (LSI) ensures an exponential contraction in KL such as

 $\mathsf{KL}(\delta_x P^n \parallel \pi) \lesssim \exp$ 

ion in 
$$\chi^2$$
 such as

$$P\left(-\frac{n}{C_{\mathsf{PI}}(\pi)}\right)\chi^{2}(\delta_{x}P^{1} \parallel \pi)$$

$$\left(-\frac{n}{C_{\mathsf{LSI}}(\pi)}\right)\mathsf{KL}(\delta_x P^1 \parallel \pi)$$

# Functional inequalities for boosting

**Q**. When can we bound sup  $TV(\delta_x P^n, \pi)$  without huge overhead?  $x \in \Omega$ 

Recall  $2 \text{ TV}^2 \leq \text{KL} \leq \log(1 + \chi^2) \leq \chi^2$ . In general,  $\mathsf{KL}(\delta_{\mathbf{y}}P \parallel \pi) = \mathsf{poly}(d)$  $\chi^2(\delta_{\mathbf{y}} P \parallel \pi) = \exp(\operatorname{poly}(d))$ 

Under (**PI**), the convergence rate would have the overhead of  $\log \chi_0^2 = \text{poly}(d)$ 

Under (LSI), the convergence rate would have the overhead of  $\log KL_0 = polylog(d)$ 

**LSI** can provide  $\mathscr{R}_{\infty}$ -guarantee only with polylog overhead!



# Annealing through Gaussians

Work with the Gaussian cooling [Cousin and Vempala'18]:

 $\rightarrow$  Need a sampler for a truncated Gaussian

- $\mu_0 \to \cdots \to \mu_i \to \mu_{i+1} \to \cdots \to \mu_k \to \pi$  with  $\mu_i = \mathsf{N}(0, \sigma_i^2 I_d)|_K$  and  $\pi = \mathsf{Unif}(K)$

Proximal sampler once again!



# Sampling from a truncated Gaussian

Proximal sampler for a truncated Gaussian  $\pi(x, y) \propto \exp\left(-V(x) - \frac{1}{2h}\|x - \frac{1}{2h}$ 

## **Algorithm**: Repeat

- 1. Sample  $y_{i+1} \sim \pi^{Y|X=x_i}(y) = N(x_i, hI_d)$
- 2. Sample  $x_{i+1} \sim \pi^{X|Y=y_{i+1}}(x) \propto \exp(-V(x) V(x))$

Q. What's (1) the convergence rate and (2) query complexity of the backward step?

$$|-y||^2$$
) with  $V(x) = \frac{1}{2\sigma^2} ||x||^2 \cdot 1_K(x)$ 

$$\frac{1}{2h} \|x - y\|^2 = \mathsf{N} \left( \frac{1}{1 + h\sigma^{-2}} y_{i+1}, \frac{h}{1 + h\sigma^{-2}} I_d \right) \Big|_{K}$$



# Uniform ergodicity of proximal sampler

**Theorem.** Under suitable choices of parameters, for any  $x \in K$ , after  $n = \widetilde{O}(qd^2C_{\text{LSI}}(\pi)\log\frac{\text{poly}(d,D)}{c})$  iterations.

**Fact 1** [Bakry-Émery].  $C_{LSI}(N(\mu, \sigma^2 I_d))$ Fact 2 [Bakry-Gentil-Ledoux]. Convex truncation doesn't increase  $C_{LSI}$ 

 $\therefore C_{\mathsf{I}}$ 

- $\mathscr{R}_{a}(\delta_{x}P^{n} \| \pi) \leq \varepsilon \text{ for } \pi = \mathsf{N}(0,\sigma^{2}I_{d})|_{K}$

$$\leq \sigma^2$$

$$_{\rm SI}(\pi) \leq \sigma^2$$



# Uniform ergodicity of proximal sampler

**Theorem.** Under suitable choices of parameters, for any  $x \in K$ ,

after 
$$n = \widetilde{O}\left(d^2\sigma^2 \log \frac{\operatorname{poly}(d,D)}{\varepsilon}\right)$$
 iteration

Boost from TV (from any start)  $\rightarrow \mathscr{R}_{\infty}$ 

- $\mathsf{TV}(\delta_x P^n, \pi) \leq \varepsilon \text{ for } \pi = \mathsf{N}(0, \sigma^2 I_d)|_{K}$ 
  - tions.

 $\mathscr{R}_{\infty}(\mu P^{n} \| \pi) \leq \varepsilon \text{ for } \pi = \mathsf{N}(0, \sigma^{2} I_{d})|_{K}$ 



# Query complexity of proximal sampler

Use a rejection sampling to implement the backward-step

- **Theorem.** [Complexity] For a well-rdd convex K, failure prob.  $\delta \in (0,1)$ , target acc.  $\varepsilon \in (0,1)$ ,  $\exists$  parameters h, N such that with probability  $\geq 1 - \delta$ ,
  - the Proximal sampler with *M*-warm start ensures  $\mathscr{R}_{\infty}(\operatorname{law}(X_n) \| \pi) \leq \varepsilon$ by using  $\widetilde{O}(Md^2\sigma^2 \operatorname{polylog} \frac{D}{\delta\varepsilon})$  membership queries in expectation.



# Annealing through Gaussians

$$\mu_0 \to \cdots \to \mu_i \to \mu_{i+1} \to \cdots \to \mu_k \to \pi$$

Set 
$$\sigma_0^2 = 1/d$$
, and update according to  
 $\sigma_{i+1}^2 \leftarrow \begin{cases} \sigma_i^2 \left(1 + \frac{1}{d}\right) & \text{if } d^{-1} \le \sigma_i^2 \le 1 \\ \sigma_i^2 \left(1 + \frac{\sigma_i^2}{d}\right) & \text{if } 1 \le \sigma_i^2 \le d \end{cases}$ 

Employ the proximal sampler within the Gaussian cooling [Cousin and Vempala'18]:  $\pi$  with  $\mu_i = N(0, \sigma_i^2 I_d) |_K$  and  $\pi = Unif(K)$ 



# **Annealing through Gaussians**

$$\sigma_{i+1}^2 \leftarrow \begin{cases} \sigma_i^2 \left(1 + \frac{1}{d}\right) & \text{if } d^{-1} \le \sigma_i^2 \le 1 \\ \sigma_i^2 \left(1 + \frac{\sigma_i^2}{d}\right) & \text{if } 1 \le \sigma_i^2 \le d \end{cases}$$

Query complexity of Gaussian sampling from an O(1)-warm start:  $d^2\sigma^2$ 

1. During  $d^{-1} \leq \sigma_i^2 \leq 1$ , needs  $\widetilde{O}(d)$  phases for doubling of  $\sigma_i^2 \to \#$  queries :  $d \cdot d^2 \sigma^2 \leq d^3$ 

2. During  $1 \le \sigma_i^2 \le d$ , needs  $\widetilde{O}(d/\sigma_i^2)$  phases for doubling  $\rightarrow \#$  queries :  $d/\sigma^2 \cdot d^2\sigma^2 \le d^3$ 

 $\therefore$  Total query complexity through annealing:  $d^3$ 



# Wrap-up in the last phase

Wrap-up for uniform sampling in the last phase  $\mu_k = \mathsf{N}(0, dI_d)|_{\kappa}$ 

Use the boosting for uniform sampling via LSI

$$_{K} \rightarrow \pi = \text{Unif}(K)$$

## $\rightarrow$ INO (or proximal sampler)'s complexity: $\widetilde{O}(d^2C_{1SI}) = \widetilde{O}(d^2D^2)$ in the last phase

# Wrap-up in the last phase

Wrap-up for uniform sampling in the last phase  $\mu_k = \mathsf{N}(0, dI_d)$ 

**Better way**?

<u>Rationale</u>: A log-concave dist. has a sub-exponential tail

$$_{K} \rightarrow \pi = \text{Unif}(K)$$

## Can work with the uniform dist. $\hat{\pi}$ over a truncated body $K \cap B_{O(d^{1/2})}(0)$ due to $\mathscr{R}_{\infty}(\hat{\pi} \parallel \pi) = O(1)$

 $P_{\pi}(\|X - \mu\| \ge t\sqrt{d}) \le \exp(-t + 1)$  when  $\mathbb{E}_{\pi}[\|X\|^2] = O(d)$ 



# Wrap-up in the last phase

Wrap-up for uniform sampling in the last phase

 $\mu_k = \mathsf{N}(0, dI_d)$ 

After truncation by  $B_{O(d^{1/2})}(0)$ :  $D = O(d^{1/2})$  so C

$$_{K} \rightarrow \pi = \text{Unif}(K)$$

$$C_{\text{LSI}}(\hat{\pi}) = O(D^2) = O(d)$$

 $\therefore$  INO's complexity is  $d^3$  for uniform sampling in the last phase



# Putting together

1. Annealing through Gaussian

 $\widetilde{O}(d^3 \operatorname{pol}$ 

2. Uniform sampling in the last phase  $\widetilde{O}(d^3 \operatorname{poly})$ 

 $\therefore \widetilde{O}(d^3 \operatorname{polylog} \frac{D}{\delta \varepsilon}) \text{ membership queries for uniform sampling with } \mathscr{R}_{\infty}\text{-guarantee}$ 

 $\rightarrow$  Matches the prior best known complexity in TV [Cousin and Vempala'18]

$$\left( \text{lylog} \frac{D}{\delta} \right)$$
 queries

$$\operatorname{ylog} \frac{D}{\delta \varepsilon}$$
) queries

