# Mitigating Externalities while Learning

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# **Example of externalities:** factories on a river



### **Example: factories on a river**

**Upstream utility:**  $\pi_1(q_1)$  increases with  $q_1$ **Downstream utility:**  $\pi_2(q_1, q_2)$  increases with  $q_2$  and decreases with  $q_1$ 

**Upstream optimal strategy:**  $q_1^{\star} = \arg \max \pi_1(q)$ 

**Downstream optimal strategy:**  $q_2^{\star} = \arg \max \pi_2(q_1^{\star}, q)$ 

Social inefficiency:  $\pi_1(q_1^{\star}) + \pi_2(q_1^{\star}, q_2^{\star}) < max_{q_1, q_2} \{ \pi(q_1) + \pi(q_1, q_2) \}$ 





## **Recovering social optimum**

**Idea:** allowing proprietary rights to restore social efficiency



 $\rightarrow$  recover social optimum

 $\max \pi_1(q_1) + \pi_2(q_1, q_2)$  $q_1, q_2$ 

### Issue: the players do not know their utilities!



### We can recover social optimum through transfers

### What if we learn utility functions over time?



#### For instance: Ressource sharing





### **Model with transfers**

At each time *t*:

- Downstream player proposes payment  $(\tau, \tilde{a}) \in \mathbb{R}_+ \times [K]$
- 2. Upstream player:
- observes  $(\tau(t), \tilde{a}_t)$
- plays  $A_t \in [K]$
- 3. Downstream player: plays  $B_t \in [K]$ 
  - gets reward  $X_{A_t,B_t}(t)$
  - observes separately



**Downstream Player** 

- receives and observes:  $Z_{A_t}(t) + 1_{A_t = \tilde{a}} \tau$ 

$$\begin{array}{l} \textbf{f} ) - \mathbf{1}_{A_t = \tilde{a}} \tau \\ \textbf{y} \ X_{A_t, B_t}(t) \text{ and } \mathbf{1}_{A_t = \tilde{a}} \end{array} \end{array}$$





**Downstream Player** 



# Designing no regret strategies

Upstream player is a no-regret learner

## **Downstream policy**

**Idea:** run batched binary searches to find  $\tau_a^{\star}$ : the minimal incentive to have  $A_t = a$ 

- Propose payment  $(a, \tau)$  for  $T^{\alpha}$  successive time steps
- Observe  $T^{\neq}$  the number of times upstream did not pull a

Using the **no-regret assumption**, w.h.p.

• If 
$$T^{\neq} > CT^{\alpha\kappa+\beta}$$
, then  $\tau_a^{\star} > \tau - \frac{1}{T^{\beta}}$   
• If  $T^{\neq} < T^{\alpha} - CT^{\alpha\kappa+\beta}$ , then  $\tau_a^{\star} < \tau + \frac{1}{T^{\beta}}$ 

For any  $\beta$  s.t.  $\alpha \kappa + \beta < \alpha$ 



### **Downstream policy**

### Theorem



 $\rightarrow$  most of the regret is due to waiting for the upstream player learning

- $\rightarrow$  upstream and downstream players can typically use UCB as a subroutine

 $\rightarrow$  the faster does the upstream player learns, the better for the downstream one

### Conclusion

### **Summary:**

- Study a repeated two player games with an upstream/downstream relation
- We propose a downstream algorithm that works for general upstream policies

#### **Direct extensions:**

- Instance dependent bounds
- Anytime policy
- Extension to linear contextual case

### Perspectives

**Higher level questions:** 

- More general interactions between multiple learning agents
- Propose black box independent strategies
- Potential long term strategic manipulations

Thank you!