Advice Querying under Budget Constraint for Online Algorithms

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Online algorithms

- Minimization (or maximization) problem
- data revealed sequentially $X = (x_1, \ldots, x_k)$
- Competitive ratio: $CR(ALG) = \sup_{\mathcal{I}} \frac{ALG(X)}{OPT(X)}$

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Learning-augmented algorithms

- predictions $Y = (y_1, \ldots, y_k)$ of $X = (x_1, \ldots, x_k)$, $\eta = d(X, Y)$
- **Robustness**: $\eta \to \infty \implies CR(ALG, Y) = O(CR(ALG))$
- Consistency: $\eta \rightarrow 0 \implies CR(ALG, Y) = O(1)$
- λ : parameter for tuning the consistency/robustness tradeoff

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- one prediction for each unknown parameter
- the predictions are given as input
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What if ...

- limited number of predictions?
- the algorithm chooses when to query them?
- guarantees on their quality?

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Ski-rental problem

- cost of renting = 1, cost of buying = b > 1
- x =unknown number of snow days

• OPT(x) = min(x, b), CR(DET) =
$$2 - \frac{1}{b}$$
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CR(RAND) = $(1 - (1 - 1/b)^b)^{-1} \rightarrow \frac{e}{e-1}$

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- Q_t prediction of (x t > b)
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Algorithm

- ALG_t Rent until t then run $ALG(p_t)$
- $CR(ALG_t) \leq \frac{t}{b} + CR(ALG(p_t))$

5/7

Secretary problem

- *n* values observed sequentially in uniformly random order
- objective: Select the maximum
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- can be queried by the algorithm up to B times

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Algorithm

- adaptive threshold: restarted 1/e-rule
- trust the recommendation if $p > p^*$
- $\mathsf{Pr}(\mathsf{ALG} \ \mathsf{succeeds}) o 1$ for $B o \infty$

6/7

Preemptive non-clairvoyant scheduling problem

- N jobs of unknown sizes to be scheduled on a single machine
- objective: minimize the sum of the completion times
- Round-robin algorithm: CR(RR) = 2

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Algorithm

- query the sizes of B random jobs
- Run concurrently OPT on the B jobs with known sizes and RR on the N-B jobs with unknown sizes

• Competitive ratio
$$2 - \frac{B(B-1)}{N(N-1)}$$