

# Accelerating Value Iteration with Anchoring

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## Value Iteration

With the given Markov Decision Process  $(\mathcal{S}, \mathcal{A}, P, r, \gamma)$  and policy  $\pi$ , define Bellman operators as

$$\begin{aligned}T^\pi V(s) &= \mathbb{E}_{a \sim \pi(\cdot | s), s' \sim P(\cdot | s, a)} [r(s, a) + \gamma V(s')], \\T^* V(s) &= \sup_{a \in \mathcal{A}} \{r(s, a) + \gamma \mathbb{E}_{s' \sim P(\cdot | s, a)} [V(s')]\}.\end{aligned}$$

Value Iteration (VI) is a fixed-point iteration with  $T (= T^\pi \text{ or } T^*)$ :

$$V^{k+1} = TV^k. \quad (\text{VI})$$

For  $0 < \gamma < 1$ , VI exhibits the rate

$$\|TV^k - V^k\|_\infty \leq (1 + \gamma)\gamma^k \|V^0 - V^*\|_\infty.$$

**Question)** If  $\gamma \approx 1$ , VI has rate  $\mathcal{O}(1)$ . Can we accelerate this in terms of Bellman error?

## Anchored Value Iteration

*Anchored Value Iteration (Anc-VI):*

$$V^k = \beta_k V^0 + (1 - \beta_k) T V^{k-1}$$

where  $\beta_k = 1 / (\sum_{i=0}^k \gamma^{-2i})$ .

We call the  $\beta_k V_0$  term the *anchor term* since it serves to pull the iterates toward the starting point  $V_0$ . The strength of the anchor mechanism diminishes since  $\beta_k$  is a decreasing sequence.

## Accelerated rate of Anc-VI

### Theorem

For  $0 < \gamma < 1$ , if  $V^0 \leq TV^0$  or  $V^0 \geq TV^0$ , Anc-VI exhibits the rate

$$\begin{aligned}\|TV^k - V^k\|_\infty &\leq \frac{(\gamma^{-1} - \gamma)(1 + \gamma - \gamma^{k+1})}{(\gamma^{k+1})^{-1} - \gamma^{k+1}} \|V^0 - V^*\|_\infty \\ &= \left( \frac{1}{k+1} + \frac{k}{k+1}\epsilon + O(\epsilon^2) \right) \|V^0 - V^*\|_\infty\end{aligned}$$

where  $\epsilon = 1 - \gamma$  and the big- $\mathcal{O}$  notation considers the limit  $\epsilon \rightarrow 0$ .

This implies that

$$\underbrace{\frac{(\gamma^{-1} - \gamma)(1 + \gamma - \gamma^{k+1})}{(\gamma^{k+1})^{-1} - \gamma^{k+1}}}_{\text{rate of Anc-VI}} < \underbrace{(1 + \gamma)\gamma^k}_{\text{rate of VI}} \quad \text{for } 0 < \gamma < 1.$$

## Complexity lower bound

### Theorem

Let  $k \geq 0$ ,  $n \geq k + 2$ ,  $0 < \gamma \leq 1$ , and  $V^0 \in \mathbb{R}^n$ . Then there exists an MDP with  $|\mathcal{S}| = n$  and  $|\mathcal{A}| = 1$  (which implies the Bellman consistency and optimality operator all coincide as  $T: \mathbb{R}^n \rightarrow \mathbb{R}^n$ ) such that  $T$  has a fixed point  $V^*$  satisfying  $V^0 \leq V^*$  and

$$\|TV^k - V^k\|_\infty \geq \frac{\gamma^k}{\sum_{i=0}^k \gamma^i}{}^3 \|V^0 - V^*\|_\infty$$

for any iterates  $\{V^i\}_{i=0}^k$  satisfying

$$V^i \in V^0 + \text{span}\{TV^0 - V^0, TV^1 - V^1, \dots, TV^{i-1} - V^{i-1}\}$$

for  $i = 1, \dots, k$ .

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<sup>3</sup>Our Theorem improves upon the prior state-of-the-art complexity lower bound established in the (V. Goyal, 2020) by a factor  $1 - \gamma^{k+1}$ .

V. Goyal and J. Grand-Clément. A first-order approach to accelerated value iteration. Operations Research, 2022.

## Optimality of Anc-VI

By previous two theorems,

$$\underbrace{\frac{\gamma^k}{\sum_{i=0}^k \gamma^i}}_{\text{lower bound}} \leq \underbrace{\frac{(\gamma^{-1} - \gamma)(1 + \gamma - \gamma^{k+1})}{(\gamma^{k+1})^{-1} - \gamma^{k+1}}}_{\text{upper bound (rate of Anc-VI)}} \leq \frac{4\gamma^k}{\underbrace{\sum_{i=0}^k \gamma^i}_{4 \times \text{lower bound}}} .$$

This implies that Anc-VI is optimal up to a constant of factor 4 while VI is suboptimal.

## Summary

Anc-VI reduces the Bellman error faster than VI. Furthermore, Anc-VI is optimal method up to factor of 4.

We also show that the anchoring mechanism provides the same benefit in the approximate VI and Gauss–Seidel VI setups as well.