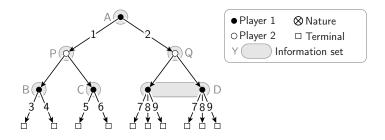
Polynomial-Time Linear-Swap Regret Minimization in Imperfect-Information Sequential Games

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Extensive-form games (EFGs)



Given a set Φ of strategy transformations $\phi : \mathcal{X} \to \mathcal{X}$, a Φ -regret minimizer is an *online learning* algorithm that minimizes

$$\Phi\text{-}\mathsf{Reg}^{(T)} \coloneqq \max_{\phi \in \Phi} \sum_{t=1}^{T} u^{(t)}(\phi(\mathbf{x}^{(t)})) - \sum_{t=1}^{T} u^{(t)}(\mathbf{x}^{(t)})$$

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Celebrated Result: If all players of the game are Φ -regret minimizers, then the empirical frequency of play converges to the set of Φ -equilibria.



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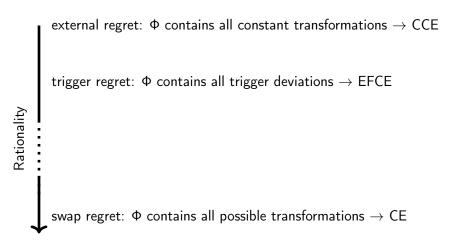
external regret: Φ contains all constant transformations \rightarrow CCE

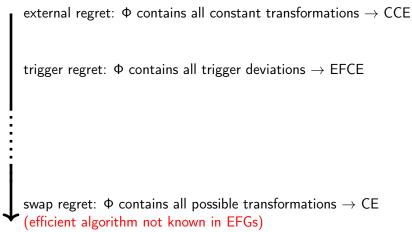


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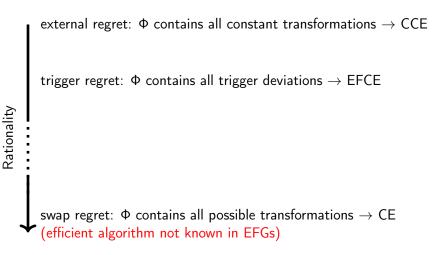
trigger regret: Φ contains all trigger deviations \rightarrow EFCE







Rationality



Big Challenge: Largest *tractable* Φ in extensive-form games?

Rationality

```
external regret: \Phi contains all constant transformations \rightarrow CCE
trigger regret: \Phi contains all trigger deviations \rightarrow EFCE
linear-swap regret (this work)
swap regret: \Phi contains all possible transformations \rightarrow CE (efficient algorithm not known in EFGs)
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Big Challenge: Largest *tractable* Φ in extensive-form games?

Let ${\mathcal Q}$ be the set of all $\mathit{reduced}$ sequence-form strategies. Then

$$\Phi \coloneqq \{ \boldsymbol{x} \mapsto \boldsymbol{A} \boldsymbol{x} : \boldsymbol{A} \in \mathbb{R}^{d \times d}, \text{ with } \boldsymbol{A} \boldsymbol{x} \in \mathcal{Q} \quad \forall \, \boldsymbol{x} \in \mathcal{Q} \}.$$

Resulting notion of regret is "linear-swap regret".

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Resulting notion of regret is "linear-swap regret".

Why care?

- in EFGs it contains the trigger deviations used for EFCE
- in NFGs it is equal to the swap deviations and gives CE
- in Bayesian games with a learner vs optimizer, it is necessary to have at least a linear-swap regret minimizer [1]

[1] Mansour, Mohri, Schneider, Sivan (2022)

 Φ -regret minimizer on action set \mathcal{X} .



External regret minimizer on action set Φ .

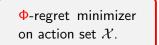
[1] Gordon, Greenwald, Marks (2008)

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 $\implies \qquad \begin{array}{c} \text{External regret mini-} \\ \text{mizer on action set } \Phi. \end{array}$

But the framework alone does *not* give us *efficient* algorithms.

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Crucial step: We characterize the set of linear-swap transformations using poly num. of constraints.

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But the framework alone does *not* give us *efficient* algorithms.

Crucial step: We characterize the set of linear-swap transformations using poly num. of constraints.

- the characterization depends on the structure of the game tree
- allows us to directly apply online convex optimization methods

Strictly between normal-form and extensive-form correlated equilibrium:

$\mathsf{CE} \subset \textbf{LCE} \subset \mathsf{EFCE}$

And, NP-hard to maximize Social Welfare, by a reduction from SAT.

- We construct a **polynomial-time** algorithm for **linear-swap regret** minimization in **EFGs**.
 - A stronger notion of sequential hindsight rationality that can be efficiently computed.
- These dynamics induce the linear-deviation correlated equilibrium.
 - It lies *strictly* between CE and EFCE.
 - Equilibrium selection is NP-hard.