# On Private and Robust Bandits

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- The agent interacts with the environment for *T* rounds.
- In each round t, the agent chooses an action  $a_t \in [K]$
- Standard reward  $r_t$  is generated independently from inlier distribution.
- After contamination, the agent observes contaminated reward  $x_t$ .

#### Definition (Finite *k*-th raw moment)

A distribution over  $\mathbb{R}$  is said to have a finite *k*-th raw moment if it is within

$$\mathcal{P}_k = \left\{ \boldsymbol{P} : \mathbb{E}_{\boldsymbol{X} \sim \boldsymbol{P}} \left[ |\boldsymbol{X}|^k 
ight] \leq 1 
ight\}, \quad k \geq 2,$$

where *k* is considered fixed but arbitrary.

#### Definition (Finite *k*-th central moment)

A distribution over  $\mathbb{R}$  is said to have a finite *k*-th central moment if it is within

$$\mathcal{P}_k^c = \left\{ \boldsymbol{P} : \mathbb{E}_{\boldsymbol{X} \sim \boldsymbol{P}} \left[ |\boldsymbol{X} - \boldsymbol{\mu}|^k \right] \leq 1 \right\}, \quad k \geq 2,$$

where  $\mu := \mathbb{E}_{X \sim P}[X] \in [-D, D]$  and  $D \ge 1$ .

## Definition (Heavy-tailed MABs with Huber contamination)

Given the corruption level  $\alpha \in [0, 1/2)$ . For each round  $t \in [T]$ , the observed reward  $x_t$  for action  $a_t$ , is sampled independently from the true distribution  $P_{a_t} \in \mathcal{P}_k$  (or  $P_{a_t} \in \mathcal{P}_k^c$ ) with probability  $1 - \alpha$ ; otherwise is sampled from some arbitrary and unknown contamination distribution  $G_{a_t} \in \mathcal{G}$ .

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## Definition (Differential Privacy for MABs)

For any  $\epsilon > 0$ , a learning algorithm  $\mathcal{M} : \mathbb{R}^T \to [K]^T$  is  $\epsilon$ -DP if for all sequences  $\mathcal{D}_T, \mathcal{D}'_T \in \mathbb{R}^T$  differing only in a single element and for all events  $E \subset [K]^T$ , we have

$$\mathbb{P}\left[\mathcal{M}(\mathcal{D}_{\mathcal{T}})\in \pmb{E}
ight]\leq \pmb{e}^{\epsilon}\cdot\mathbb{P}\left[\mathcal{M}\left(\mathcal{D}_{\mathcal{T}}'
ight)\in \pmb{E}
ight].$$

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# Regrets

- $\mu_a$ : the mean of the inlier distribution of arm  $a \in [K]$ ;
- $\mu^* = \max_{a \in [K]} \mu_a;$
- $\Pi^{\epsilon}$ : the set of all  $\epsilon$ -DP MAB algorithms;
- $\mathcal{E}_{\alpha,k}$ : the set of all instances of heavy-tailed MABs (with parameter *k*) with Huber contamination (of level  $\alpha$ ).

#### **Definition (Clean Regret)**

Fix an algorithm  $\pi \in \Pi^{\epsilon}$  and an instance  $\nu \in \mathcal{E}_{\alpha,k}$ . Then, the clean regret of  $\pi$  under  $\nu$  is given by  $\mathcal{R}_T(\pi, \nu) := \mathbb{E}_{\pi,\nu}[T\mu^* - \sum_{t=1}^T \mu_{a_t}].$ 

To capture the intrinsic difficulty of the private and robust MAB problem, we are also interested in its minimax regret.

## Definition (Minimax Regret)

The minimax regret of our private and robust MAB problem is defined as  $\mathcal{R}_{\epsilon,\alpha,k}^{\min \max} := \inf_{\pi \in \Pi^{\epsilon}} \sup_{\nu \in \mathcal{E}_{\alpha,k}} \mathbb{E}_{\pi,\nu}[T\mu^* - \sum_{t=1}^{T} \mu_{a_t}].$ 

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#### Theorem

Consider a private and robust MAB problem where inlier distributions have finite k-th raw (or central) moments ( $k \ge 2$ ). Then, its minimax regret satisifes

$$\mathcal{R}_{\epsilon,\alpha,k}^{\min(max)} = \Omega\left(\sqrt{KT} + \left(\frac{K}{\epsilon}\right)^{1-\frac{1}{k}} T^{\frac{1}{k}} + T\alpha^{1-\frac{1}{k}}\right).$$

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#### A Meta Algorithm

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# A Meta Algorithm

#### Algorithm 1 Private and Robust Arm Elimination

- 1: **Input:** Number of arms K, time horizon T, privacy budget  $\epsilon$ , Huber parameter  $\alpha \in (0, 1/2)$ , error probability  $\delta \in (0, 1]$ , inliner distribution parameters i.e., k and optional D.
- 2: Initialize:  $\tau = 0$ , active set of arms  $S = \{1, \dots, K\}$ .
- 3: for batch  $\tau = 1, 2, \ldots$  do
- 4: Set batch size  $B_{\tau} = 2^{\tau}$ .
- 5: **if**  $B_{\tau} < \mathcal{T}$  **then**
- 6: Randomly select an action  $a \in [K]$ .
- 7: Play action a for  $B_{\tau}$  times.

#### 8: else

- 9: for each active arm  $a \in S$  do
- 10: **for** *i* from 1 to  $B_{\tau}$  **do**
- 11: Pull arm a, observe contaminated reward  $x_i^a$ .
- 12: If total number of pulls reaches T, exit.
- 13: end for
- 14: Set truncation threshold  $M_{\tau}$ .
- 15: Set additional parameters  $\Phi$ .
- 16: Compute estimate  $\tilde{\mu}_a = \text{PRM}(\{x_i^a\}_{i=1}^{B_\tau}, M_\tau, \Phi).$
- 17: end for
- 18: Set confidence radius  $\beta_{\tau}$ .
- 19: Let  $\widetilde{\mu}_{\max} = \max_{a \in S} \widetilde{\mu}_a$ .
- 20: Remove all arms a from S s.t.  $\tilde{\mu}_{max} \tilde{\mu}_a > 2\beta_{\tau}$ .
- 21: end if
- 22: end for

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#### Algorithm 2 PRM for the finite raw moment case

- 1: Input: A collection of data  $\{x_i\}_{i=1}^n$ , truncation parameter M, additional parameters  $\Phi = \{\epsilon\}$ .
- 2: for i = 1, 2, ..., n do

3: Truncate data 
$$\bar{x}_i = x_i \cdot \mathbb{1}_{\{|x_i| \le M\}}$$

- 4: end for
- 5: Return private estimate  $\widetilde{\mu} = \frac{\sum_{i=1}^{n} \overline{x}_i}{n} + \operatorname{Lap}(\frac{2M}{n\epsilon}).$

#### Theorem (Performance Guarantees)

Consider a private and robust MAB with inlier distributions satisfying Definition 1 and  $0 < \alpha \le \alpha_1 \in (0, 1/2)$ . Let Algorithm 1 be instantiated with Algorithm 2. Set  $\mathcal{T} = \Omega(\frac{\log(1/\delta)}{\alpha_1})$  and  $\delta = 1/T$ . Then Algorithm 1 is  $\epsilon$ -DP with its regret upper bound

$$\mathcal{R}_T = O\left(\sqrt{KT\log T} + \left(\frac{K\log T}{\epsilon}\right)^{\frac{k-1}{k}} T^{\frac{1}{k}} + T\alpha_1^{1-\frac{1}{k}} + \frac{K\log T}{\alpha_1}\right).$$

Algorithm 3 PRM for the finite central moment case

- 1: Input: A collection of data  $\{x_i\}_{i=1}^{2n}$ , truncation parameter M, additional parameters  $\Phi = \{\epsilon, D, r\}, r \in \mathbb{R}$ .
- 2: // First step: initial estimate
- 3:  $B_j = [j, j+r), j \in \mathcal{J} = \{-D, -D+r, \dots, D-r\}.$
- 4: Compute private histogram using the first fold of data:  $\widetilde{p}_j = \frac{\sum_{i=1}^n \mathbb{1}_{\{X_i \in B_j\}}}{n} + \operatorname{Lap}\left(\frac{2}{n\epsilon}\right)$
- 5: Get the initial estimate  $J = \arg \max_{j \in \mathcal{J}} \widetilde{p}_j$ .
- 6: // Second step: final estimate
- 7: Get final estimator using the second fold of data:  $\widetilde{\mu} = J + \frac{1}{n} \sum_{i=n+1}^{2n} (X_i J) \mathbb{1}_{\{|X_i J| \le M\}} + Lap(\frac{2M}{n\epsilon}).$

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#### Theorem (Performance Guarantees, $\alpha = 0$ )

Let Algorithm 1 be instantiated with Algorithm 3. Set  $T = \Omega(\frac{\log(D/\delta)}{\epsilon})$  and  $\delta = 1/T$ . Then, Algorithm 1 is  $\epsilon$ -DP with its regret upper bound

$$\mathcal{R}_{T} = O(\sqrt{KT \log T} + (K \log T/\epsilon)^{\frac{k-1}{k}} T^{\frac{1}{k}} + \gamma),$$

where  $\gamma := O(KD \log(DT)/\epsilon)$ .

#### Theorem (Performance Guarantees, $\alpha > 0$ )

For  $\alpha \leq \alpha_1 \in (0, 0.133)$ , let Algorithm 1 be instantiated with Algorithm 3. Set  $\delta = 1/T$ , then Algorithm 1 is  $\epsilon$ -DP with its regret upper bound

$$\mathcal{R}_T = O(\sqrt{KT\log T} + (K\log T/\epsilon)^{\frac{k-1}{k}} T^{\frac{1}{k}} + T\alpha_1^{1-\frac{1}{k}} + \hat{\gamma}),$$

where  $\hat{\gamma} := O\left(\frac{DK\log T}{\alpha_1^2} + \frac{\iota DK\log T}{\epsilon} + \frac{DK\log(DT)}{\epsilon}\right)$  and  $\iota = \frac{1-\alpha}{0.249-\alpha}$ .

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#### 6 Experiments

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## Experiments

- PRAE-R: Our Algorithm for Finite Raw Moment Case
- PRAE-C: Our Algorithm for Finite Central Moment Case
- DPRSE [Tao et al., 2021]: DP heavy-tailed MAB
- RUCB [Kapoor et al., 2019]: non-private robust algorithm

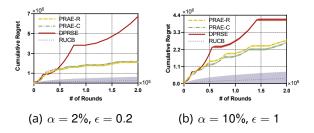


Figure: Experimental results under Pareto distribution

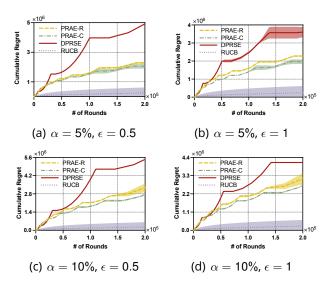


Figure: Experimental results under Student's t reward

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# Thank you!

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