Exact Privacy Guarantees for Markov Chain Implementations of the Exponential Mechanism

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- We propose an exact finite-runtime algorithm for implementing the exponential mechanism with exact ε-DP guarantees using atomic regeneration.
- We show that our proposed algorithm relies on the confidential data to avoid the worst-case mixing times found in distance convergence arguments.
- We derive two modifications of the previous algorithm that demonstrate a new three way trade-off between, privacy, utility, and runtime.

Notation:

 $\begin{cases} \mathcal{X} \triangleq: \text{sample space of 1 individual's data} \\ (\mathcal{Y}, \mathcal{F}) \triangleq \text{output space} \\ \mathcal{M} \triangleq \{\mu_X \mid X \in \mathcal{X}^n\} \text{ release mechanism} \end{cases}$

Differential privacy (Dwork et al, 2006)

A mechanism \mathcal{M} satisfies (ϵ, δ) -DP if, for all $B \in \mathcal{F}$ and adjacent $X, X' \in \mathcal{X}^n$:

 $\mu_X(B) \leq e^{\epsilon} \mu_{X'}(B) + \delta.$

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When $\delta = 0$, we say a mechanism satisfies ϵ -DP.

Private selection and the exponential mechanism

Private selection:

• Goal: minimize loss function L_X while satisfying ϵ -DP

$$L_X: \mathcal{X}^n \times \mathcal{Y} \mapsto [0, \infty]$$

• Key ingredient: bounded sensitivity of L_X . For all adjacent $X, X' \in \mathcal{X}^n$:

$$|L_X(y) - L_{X'}(y)| \le \Delta_L < \infty.$$

Exponential mechanism (McSherry and Talwar, 2007)

A sample from density f_X with the form:

$$f_X(y) \propto \exp\left(-rac{\epsilon L_X(y)}{2\Delta_L}
ight),$$

with respect to a common base measure $\nu(y)$ over $(\mathcal{Y}, \mathcal{F})$ satisfies ϵ -DP.

Many generic ϵ -DP algorithms are not exactly implementable!

Why can't we just use MCMC?

- MCMC approximation has a privacy cost
- Heuristic MCMC convergence measures tell us nothing about said cost

Approximation δ cost (Li et al, 2016)

A sequence of mechanisms $\mathcal{M}_m \triangleq \{\mu_{m,X} \mid x \in \mathcal{X}\}$ approximating the exponential mechanism, \mathcal{M}_m , as $m \ge \tau(\alpha)$ is $(\epsilon, \delta_\alpha)$ -DP where $\delta_\alpha \triangleq \alpha(1 + e^{\epsilon})$ if

$$\tau(\alpha) \triangleq \sup_{X \in \mathcal{X}^n} \inf\{t \ge 0 \mid \|\mu_{t,X} - \mu_X\|_{\mathrm{TV}} \le \alpha\}.$$

Current approach: bounding distributional distances between the MCMC approximation and the target distribution (ex: Ganesh and Talwar, 2020)

Problems with existing approaches:

- Asymptotic rates can't be used to calculate finite-chain privacy loss
- Need to bound distances for worst-case slowest mixing chains
- Methods don't exactly satisfy ϵ -DP

For a Markov chain with stationary distribution μ_X and transition kernel Π_X :

- Associate with each state a binary indicator $\rho \in \{0, 1\}$ indicating regeneration (i.e. return to the same state)
- Let $\{\tau\}_{t=1}^{\infty}$ be the sequence of **regeneration times** for the MC (i.e. time between states when $\rho = 1$)

Each τ_t is IID (can drop t index) $\implies \mu_X$ has an infinite mixture form:

$$\mu_X(A) = \sum_{m=1}^{\infty} \frac{\mathbb{P}(\tau \ge m)}{\mathbb{E}[\tau]} \, \mathbb{P}(Y_m \in A \mid \tau \ge m).$$

(Lee et al, 2014) show that if the regeneration state is a **singleton atom**, then we can sample from μ_X using Bernoulli factories (Huber, 2013).

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Our proposal: confidential artificial atoms

Implementation choices specific to DP:

 Choose an *artifical* confidential atom *a* ∈ 𝒱 from the set of confidential results (i.e. what we would release without privacy preservation)

$$a\in rginf_{y\in\mathcal{Y}} L_X(y).$$

• First sample from $\tilde{\mu}_X$, where:

$$\tilde{\mu}_X = (1-k)\mu_X + k\xi_a,$$

then condition on $Y \neq a$ to sample from μ_X .

 Assumptions about the state space (such as compact Xⁿ) help to satisfy our privacy AND our sampling assumptions

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• Many different possible choices for chain modification (ex: Brockwell and Kadane, 2005)

Theorem: modified Metropolis-Hastings perfect sampling for privacy

Let Π_X be the transition kernel for a Metropolis-Hastings Markov Chain with symmetric proposals q. We can construct a Markov chain on the extended space with proposals:

$$\tilde{q}(y,y') = \frac{1}{2} \left[q_X(y,y') + \mathbb{1}_{\{y'=a\}} \right],$$

And an algorithm to sample from density f_X that satisfies ϵ -DP with expected number of total proposed samples N_{prop} :

$$\mathbb{E}[N_{\text{prop}}] \leq \frac{48}{k^2(1-k)^2 \inf_{y \in \mathcal{Y}} p_{\text{Accept}}(y)},$$

where:

$$p_{\mathrm{Accept}}(y) \triangleq \int_{\mathcal{Y}} q_X(y,y') \min\left\{1, \frac{f_X(y')}{f_X(y)}\right\} d\nu(y').$$

9/15

Key property

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MCMC methods require accounting for the **slowest mixing** chain, but our method can be much faster because **the runtime depends on the realized confidential data**

Illustrative example: Laplace mechanism $(L_X(y) = \|\overline{X} - y\|_1)$ with data bounded in $[0,1]^d$

- Two original Markov chains: Metropolis-Hastings (MH) with independent uniform proposals and symmetric Laplace proposals with scale α
- \bullet Closed form expressions for worst-case δ with MH MCMC (Mengersen and Tweedie, 1996)

$$\|\mu_{X,m} - \mu_X\|_{TV} \le (1 - \beta_{\text{MCMC}})^m, \tag{1}$$

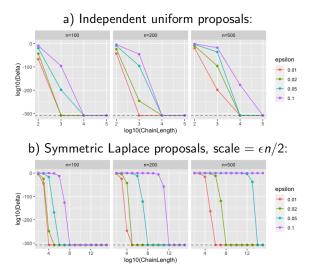
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$$\begin{cases} \beta_{\text{MCMC,Unif}} \triangleq \left(\frac{2d}{\epsilon n} (1 - e^{-\epsilon n/2d})\right)^d \\ \beta_{\text{MCMC,Lap}} \triangleq (2\alpha)^d \exp\left(-\left(\alpha d + \frac{\epsilon n}{2}\right)\right) \left(\frac{1}{\alpha} (1 - e^{-\alpha})\right)^d \end{cases}$$

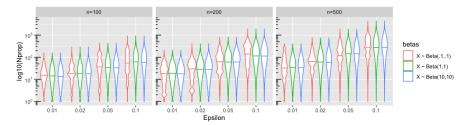
Example: *d*-dimensional Laplace mechanism in hypercube

(Dashed line = 64-bit double precision threshold)



Example: effect of central concentration

Continued example, but now with $X \sim \text{Beta}(\theta, \theta)$ with $\theta \in \{.1, 1, 10\}$ (Laplace proposals), demonstrating dependence on X:



Discussion

- Benefits
 - Satisfies ϵ -DP
 - Runtime depends on realized confidential data, and *not* the confidential data for the slowest-mixing Markov Chain
 - Only requires minorizing bound, and not properties of L_X (i.e. convexity, Lipschitz, etc.)
- Limitations
 - Uniform ergodicity assumption: methods do not have finite expected runtime for unbounded state spaces, like \mathbb{R}^d .
 - Minorizing constant suffers from curse of dimensionality
 - Side-channel vulnerability: multiple replications of similar queries could leak information about confidential data through runtime

Traditional analysis: privacy vs. utility

Extensions of our work: privacy vs. utility vs. runtime

- Trading off utility and runtime:
 - Exponential mechanisms can be implemented exactly over enumerable discrete state spaces
 - \implies corollary: if we release a sample from a discrete approximation w.p. k, then we reduce runtime at the cost of some utility
- Trading off privacy and runtime:
 - (Awan and Rao, 2021) consider rejection sampling where $\textit{N}_{\rm prop}$ is known and can leak information
 - \implies corollary: with longer artificial runtime, can release $\tilde{N}_{\text{prop}} \perp X$ with 0-DP so that $(Y, \tilde{N}_{\text{prop}})$ is ϵ -DP

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