# Newton-LESS: Sparsification without Trade-offs for the Sketched Newton Update 

Michał Dereziński* Jonathan Lacotte ${ }^{\dagger}$ Mert Pilanci ${ }^{\dagger}$ and Michael Mahoney ${ }^{\ddagger}$

*University of Michigan, ${ }^{\dagger}$ Stanford University, $\ddagger$ UC Berkeley NeurIPS 2021


## Newton's method in composite optimization

Find: $\quad \mathbf{x}^{*}=\underset{\mathbf{x}}{\operatorname{argmin}} f(\mathbf{x}), \quad$ for $\quad f(\mathbf{x})=\sum_{i=1}^{n} f_{i}(\mathbf{x})$


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Newton step: $\quad \mathbf{p}_{t}=[\underbrace{\nabla^{2} f\left(\mathbf{x}_{t}\right)}_{d \times d \text { Hessian } \mathbf{H}}]^{-1} \underbrace{\nabla f\left(\mathbf{x}_{t}\right)}_{d \times 1 \text { gradient }}$


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Newton estimate: $\widehat{\mathbf{p}}_{t}=[\underbrace{\nabla^{2} \hat{f}\left(\mathbf{x}_{t}\right)}]^{-1} \underbrace{\nabla f\left(\mathbf{x}_{t}\right)}$

$$
\text { Hessian estimate } \widehat{\mathbf{H}} \quad d \times 1 \text { gradient }
$$



## Computing the Hessian

$$
\nabla^{2} f(\mathbf{x})=\sum_{i=1}^{n} \nabla^{2} f_{i}(\mathbf{x})=\overbrace{\mathbf{A}_{f}(\mathbf{x})^{\top} \mathbf{A}_{f}(\mathbf{x})}^{\text {Cost: } O\left(n d^{2}\right)}
$$

Example: Generalized Linear Model

$$
\begin{aligned}
f(\mathbf{x}) & =\frac{1}{n} \sum_{i=1}^{n} \ell_{i}\left(\phi_{i}^{\top} \mathbf{x}\right) \\
\nabla^{2} f(\mathbf{x}) & =\frac{1}{n} \sum_{i=1}^{n} \ell_{i}^{\prime \prime}\left(\phi_{i}^{\top} \mathbf{x}\right) \phi_{i} \phi_{i}^{\top}
\end{aligned}
$$



## Newton Sketch [PW17]

$$
\widetilde{\mathbf{x}}_{t+1}=\widetilde{\mathbf{x}}_{t}-\mu_{t}(\overbrace{\mathbf{A}_{f}\left(\widetilde{\mathbf{x}}_{t}\right)^{\top} \mathbf{S}_{t}^{\top} \mathbf{S}_{t} \mathbf{A}_{f}\left(\widetilde{\mathbf{x}}_{t}\right)}^{\tilde{\mathbf{A}}_{t}^{\top} \tilde{\mathbf{A}}_{t} \approx \nabla^{2} f\left(\mathbf{x}_{t}\right)})^{-1} \nabla f\left(\widetilde{\mathbf{x}}_{t}\right)
$$



## Example 1: Gaussian Newton Sketch

Sketching matrix $\mathbf{S}_{t}$ has i.i.d. Gaussian entries
Pros

- Strong convergence
- Robust to the worst case


Extension: Sub-gaussian embeddings, e.g., with i.i.d. random sign entries

## Example 2: Sub-Sampled Newton

Randomly select $m$ rows of $\mathbf{A}_{f}\left(\mathbf{x}_{t}\right)$

Pros

- Computationally cheap


## Cons

- Weaker convergence
- Sensitive to the worst case


Extension: Importance sampling, e.g., according to leverage scores

## LESS Embeddings: Fast Gaussian-like Sketches

LEverage Score Sparsified (LESS) Embeddings:
Leverage Score Sampling + Sparse Embedding Matrices


Introduced by [DLDM21] "Sparse sketches with small inversion bias", COLT'21.

## Newton-LESS: Sparsity without trade-offs


$s$ non-zeros per row


Convergence Rate $=\left(\mathbb{E} \frac{\left\|\Delta_{T}\right\|^{2}}{\left\|\Delta_{0}\right\|^{2}}\right)^{1 / T}$ where $\Delta_{t}=\widetilde{\mathbf{x}}_{t}-\mathrm{x}^{*}$


## Same plot on real data


(a) High-coherence synthetic

(c) CIFAR-10 dataset

(b) Musk dataset

(d) WESAD dataset

## Main result: Problem-independent local convergence

Assumptions: Hessian $\mathbf{H}=\nabla^{2} f\left(\mathbf{x}^{*}\right)$ is positive definite and $f$ is (a) self-concordant, or (b) has a Lipschitz continuous Hessian.

Sketching matrix: Gaussian, sub-Gaussian, or LESS embedding with sketch size $m$ at least $C d \log (d T / \delta)$

## Theorem

There is a neighboorhood $U$ containing $\mathbf{x}^{*}$ such that if $\widetilde{\mathbf{x}}_{0} \in U$, then we can choose step size $\mu_{t}$ so that:

$$
\left(\mathbb{E}_{\delta} \frac{\left\|\Delta_{T}\right\|_{\mathbf{H}}^{2}}{\left\|\Delta_{0}\right\|_{\mathbf{H}}^{2}}\right)^{1 / T} \approx_{\epsilon} \frac{d}{m} \quad \text { for } \quad \epsilon=O\left(\frac{1}{\sqrt{d}}\right)
$$

[^0]
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$$

[^2]
## Main result: Discussion

- Same problem-independent $\left(\frac{d}{m}\right)^{T}$ convergence rate for LESS and Gaussian (down to lower order terms)
- Simple analytic expression for the optimal step size $\mu_{t}$ :

$$
\widetilde{\mathbf{x}}_{t+1}=\widetilde{\mathbf{x}}_{t}-\underbrace{\left(1-\frac{d}{m}\right)}_{\mu_{t}} \widehat{\mathbf{p}}_{t}, \quad \text { when } \mathbb{E}\left[\widehat{\mathbf{p}}_{t}\right] \approx \mathbf{p}_{t}
$$

- Extension to regularized objectives $f(\mathbf{x})=f_{0}(\mathbf{x})+g(\mathbf{x})$ : the convergence rate becomes dimension-independent,

$$
\left(\mathbb{E}_{\delta} \frac{\left\|\Delta_{T}\right\|_{\mathbf{H}}^{2}}{\left\|\Delta_{0}\right\|_{\mathbf{H}}^{2}}\right)^{1 / T} \leq_{\epsilon} \frac{d_{\mathrm{eff}}}{m} \quad \text { for } \quad d_{\mathrm{eff}}=\operatorname{tr}\left(\nabla^{2} f_{0}\left(\mathbf{x}^{*}\right) \nabla^{2} f\left(\mathbf{x}^{*}\right)^{-1}\right)
$$

## Comparison to prior work

- Under quadratic objectives $f(\mathbf{x})=\|\mathbf{A x}-\mathbf{b}\|^{2}$, the convergence rate $\left(\frac{d}{m}\right)^{T}$ was previously shown only for:
(1) strictly Gaussian embeddings [LP19],
(2) Subsampled Randomized Hadamard Transform (SRHT) in a high-dimensional asymptotic limit [LLDP20].
- For general objectives and fast sketching methods, e.g.:
(1) Row sampling (Leverage Scores) [DMM06],
(2) Sparse sketches (CountSketch and SJLT) [CW17],
(3) Trigonometric sketches (SRHT and SRTT) [AC09],
the best known rate is $\left(C \log (d T / \delta) \cdot \frac{d}{m}\right)^{T}$ [PW17].
Note: Extra constant and logarithmic factors in the bound means no analytic expressions for the optimal step size $\mu_{t}$


## Analysis: Two approaches

(1) Subspace embedding

- Standard approximation guarantee for sketching methods
- Leads to subopotimal convergence rates: $\left(C \log (d T / \delta) \cdot \frac{d}{m}\right)^{T}$

$$
\mathbf{A}_{f}\left(\widetilde{\mathbf{x}}_{t}\right)^{\top} \mathbf{S}_{t}^{\top} \mathbf{S}_{t} \mathbf{A}_{f}\left(\widetilde{\mathbf{x}}_{t}\right) \approx_{\eta} \nabla^{2} f\left(\widetilde{\mathbf{x}}_{t}\right)
$$

(2) Method of inverse moments

- Originally proposed for quadratic objectives [LP19]
- Leads to precise convergence rates and optimal step sizes
- Requires inverse moments of the sketched Hessian

$$
\mathbb{E}\left[\left(\mathbf{A}_{f}\left(\widetilde{\mathbf{x}}_{t}\right)^{\top} \mathbf{S}_{t}^{\top} \mathbf{S}_{t} \mathbf{A}_{f}\left(\widetilde{\mathbf{x}}_{t}\right)\right)^{-k}\right] \quad \text { for } \quad k=1,2
$$

## Comparison of sketching methods

(1) Subspace embedding
(2) Method of inverse moments

## Sub-Gaussian Embedding

Sketching S

| $\mathbf{s}_{i}^{\top}$ |
| :---: |
|  |

Data A leverage scores

$i$-th leverage score: $\quad \ell_{i}(\mathbf{A})=i$-th diagonal entry of $\mathbf{A} \mathbf{A}^{\dagger}$

## Comparison of sketching methods

(1) Subspace embedding
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## Leverage Score Sampling [DMM06]


$i$-th leverage score: $\quad \ell_{i}(\mathbf{A})=i$-th diagonal entry of $\mathbf{A} \mathbf{A}^{\dagger}$

## Comparison of sketching methods

(1) Subspace embedding
(2) Method of inverse moments

## Uniform Sparsification [CW13]


$i$-th leverage score: $\quad \ell_{i}(\mathbf{A})=i$-th diagonal entry of $\mathbf{A} \mathbf{A}^{\dagger}$

## Comparison of sketching methods

(1) Subspace embedding
(2) Method of inverse moments

## Leverage Score Sparsification [DLDM21]

Sketching S

$i$-th leverage score: $\quad \ell_{i}(\mathbf{A})=i$-th diagonal entry of $\mathbf{A} \mathbf{A}^{\dagger}$

## Implementing LESS Embeddings

(1) Worst-case implementation (LESS)

- Preprocessing cost: $O\left(\operatorname{nnz}(\mathbf{A}) \log n+d^{3} \log d\right)$ Approximating leverage scores $\ell_{i}(\mathbf{A})$ [DMIMW12]
- Sketching cost: $O\left(m d^{2}\right)$

Sparse matrix multiplication SA

$$
\mathrm{Cost}=O\left(\operatorname{nnz}(\mathbf{A}) \log n+m d^{2}\right)
$$

(2) Practical implementation (LESS-uniform)

- Use a uniformly sparsified sketch with $\alpha d$ non-zeros per row
- If $\alpha \geq \frac{n}{d} \max _{j} \ell_{j}(\mathbf{A})$, then we recover theoretical guarantees

$$
\text { Cost }=O\left(\alpha m d^{2}\right)
$$

$\mathrm{nnz}(\mathbf{A})=$ number of non-zeros in matrix $\mathbf{A}$.

## Experiments: Quadratic objective


(a) High-coherence synthetic matrix


(b) WESAD dataset

We use sketch size $m=4 d$, and LESS-uniform has $d$ non-zeros per row.

## Experiments: Logistic regression


(a) WESAD dataset


(b) CIFAR-10 dataset


$$
\min _{\mathbf{x} \in \mathbb{R}^{d}} \frac{1}{n} \sum_{i=1}^{n} \log \left(1+\exp \left(-b_{i} \mathbf{a}_{i}^{\top} \mathbf{x}\right)\right)+\frac{\lambda}{2}\|\mathbf{x}\|_{2}^{2}
$$

We use sketch size $m=d / 2$. Bottom plots report the CPU and GPU wall-clock times to reach a $10^{-6}$ accurate solution.

## Conclusions

- Newton-LESS: Sparsification without trade-offs
(1) Per-iteration efficiency of Sub-Sampled Newton
(2) Same convergence rate as Gaussian Newton Sketch
- Sparse sketching can beat Sub-Sampling...
(1) ...in real-world optimization tasks
(2) ...on a variety of hardware platforms
- LESS Embeddings: Fast Gaussian-like sketches
(1) Correcting the bias in distributed optimization [DLDM21]
(2) Precise convergence rates and optimal step sizes (this work)

Code available at: https://github.com/lessketching/newtonsketch

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[^0]:    $\mathbb{E}_{\delta}$ is expectation conditioned on a $1-\delta$ probability event; $\|\mathbf{v}\|_{\mathbf{H}}=\sqrt{\mathbf{v}^{\top} \mathbf{H v}} ; \quad a \approx_{\epsilon} b$ means that $(1-\epsilon) \cdot b \leq a \leq(1+\epsilon) \cdot b$

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