# Newton-LESS: Sparsification without Trade-offs for the Sketched Newton Update

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### Newton's method in composite optimization

Find: 
$$\mathbf{x}^* = \underset{\mathbf{x}}{\operatorname{argmin}} f(\mathbf{x}), \quad \text{for} \quad f(\mathbf{x}) = \sum_{i=1}^n f_i(\mathbf{x})$$

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Newton step: 
$$\mathbf{p}_t = \begin{bmatrix} \nabla^2 f(\mathbf{x}_t) \end{bmatrix}^{-1} \underbrace{\nabla f(\mathbf{x}_t)}_{d \times d \text{ Hessian } \mathbf{H}} \overset{1}{d \times 1 \text{ gradient}}$$



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Newton estimate:  $\widehat{\mathbf{p}}_t = \begin{bmatrix} \nabla^2 \widehat{f}(\mathbf{x}_t) \end{bmatrix}^{-1} \underbrace{\nabla f(\mathbf{x}_t)}_{d \times 1 \text{ gradient}}$ 



# Computing the Hessian

$$\nabla^2 f(\mathbf{x}) = \sum_{i=1}^n \nabla^2 f_i(\mathbf{x}) = \overbrace{\mathbf{A}_f(\mathbf{x})^\top \mathbf{A}_f(\mathbf{x})}^{\text{Cost: }O(nd^2)}$$

Example: Generalized Linear Model

$$\begin{split} f(\mathbf{x}) &= \frac{1}{n} \sum_{i=1}^{n} \ell_i(\phi_i^{\mathsf{T}} \mathbf{x}), \\ \nabla^2 f(\mathbf{x}) &= \frac{1}{n} \sum_{i=1}^{n} \ell_i''(\phi_i^{\mathsf{T}} \mathbf{x}) \phi_i \phi_i^{\mathsf{T}} \end{split}$$



## Newton Sketch [PW17]

$$\widetilde{\mathbf{x}}_{t+1} = \widetilde{\mathbf{x}}_t - \mu_t \Big( \underbrace{\widetilde{\mathbf{A}}_t^\top \widetilde{\mathbf{A}}_t \approx \nabla^2 f(\mathbf{x}_t)}_{\widetilde{\mathbf{A}}_f(\widetilde{\mathbf{x}}_t)^\top \mathbf{S}_t^\top \mathbf{S}_t \mathbf{A}_f(\widetilde{\mathbf{x}}_t)} \Big)^{-1} \nabla f(\widetilde{\mathbf{x}}_t)$$



## Example 1: Gaussian Newton Sketch

Sketching matrix  $\mathbf{S}_t$  has i.i.d. Gaussian entries

### $\mathbf{Pros}$

### Cons

• Computationally expensive

- Strong convergence
- Robust to the worst case



Extension: Sub-gaussian embeddings, e.g., with i.i.d. random sign entries

## Example 2: Sub-Sampled Newton

Randomly select m rows of  $\mathbf{A}_f(\mathbf{x}_t)$ 

### $\mathbf{Pros}$

• Computationally cheap

### Cons

- Weaker convergence
- Sensitive to the worst case



Extension: Importance sampling, e.g., according to leverage scores

### LESS Embeddings: Fast Gaussian-like Sketches

LEverage Score Sparsified (LESS) Embeddings:

Leverage Score Sampling + Sparse Embedding Matrices



Introduced by [DLDM21] "Sparse sketches with small inversion bias", COLT'21.

### Newton-LESS: Sparsity without trade-offs



Convergence Rate = 
$$\left(\mathbb{E} \frac{\|\Delta_T\|^2}{\|\Delta_0\|^2}\right)^{1/T}$$
 where  $\Delta_t = \widetilde{\mathbf{x}}_t - \mathbf{x}^*$   
Computational Cost =  $\underbrace{O(mds)}_{\text{sketch}} + \underbrace{O(md^2)}_{\text{Hessian}} + \underbrace{O(nd)}_{\text{gradient}}$ 

### Same plot on real data



### Main result: Problem-independent local convergence

Assumptions: Hessian  $\mathbf{H} = \nabla^2 f(\mathbf{x}^*)$  is positive definite and f is (a) self-concordant, or (b) has a Lipschitz continuous Hessian.

<u>Sketching matrix</u>: Gaussian, sub-Gaussian, or LESS embedding with sketch size m at least  $Cd \log(dT/\delta)$ 

### Theorem

There is a neighboorhood U containing  $\mathbf{x}^*$  such that if  $\widetilde{\mathbf{x}}_0 \in U$ , then we can choose step size  $\mu_t$  so that:

$$\left(\mathbb{E}_{\delta} \frac{\|\Delta_T\|_{\mathbf{H}}^2}{\|\Delta_0\|_{\mathbf{H}}^2}\right)^{1/T} \approx_{\epsilon} \frac{d}{m} \qquad for \quad \epsilon = O\left(\frac{1}{\sqrt{d}}\right)$$

$$\begin{split} \mathbb{E}_{\delta} \text{ is expectation conditioned on a } 1-\delta \text{ probability event;} \\ \|\mathbf{v}\|_{\mathbf{H}} = \sqrt{\mathbf{v}^{\top}\mathbf{H}\mathbf{v}}; \qquad a \approx_{\epsilon} b \text{ means that } (1-\epsilon) \cdot b \leq a \leq (1+\epsilon) \cdot b \end{split}$$

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### Main result: Discussion

- Same problem-independent  $(\frac{d}{m})^T$  convergence rate for LESS and Gaussian (down to lower order terms)
- Simple analytic expression for the optimal step size  $\mu_t$ :

$$\widetilde{\mathbf{x}}_{t+1} = \widetilde{\mathbf{x}}_t - \underbrace{(1 - \frac{d}{m})}_{\mu_t} \widehat{\mathbf{p}}_t, \text{ when } \mathbb{E}[\widehat{\mathbf{p}}_t] \approx \mathbf{p}_t.$$

• Extension to regularized objectives  $f(\mathbf{x}) = f_0(\mathbf{x}) + g(\mathbf{x})$ : the convergence rate becomes *dimension-independent*,

$$\left(\mathbb{E}_{\delta} \frac{\|\Delta_T\|_{\mathbf{H}}^2}{\|\Delta_0\|_{\mathbf{H}}^2}\right)^{1/T} \leq_{\epsilon} \frac{d_{\text{eff}}}{m} \quad \text{for} \quad d_{\text{eff}} = \operatorname{tr} \left(\nabla^2 f_0(\mathbf{x}^*) \nabla^2 f(\mathbf{x}^*)^{-1}\right)$$

## Comparison to prior work

• Under quadratic objectives  $f(\mathbf{x}) = \|\mathbf{A}\mathbf{x} - \mathbf{b}\|^2$ , the convergence rate  $(\frac{d}{m})^T$  was previously shown only for:

strictly Gaussian embeddings [LP19],

- ② Subsampled Randomized Hadamard Transform (SRHT) in a high-dimensional asymptotic limit [LLDP20].
- For general objectives and fast sketching methods, e.g.:
  - Row sampling (Leverage Scores) [DMM06],
  - Sparse sketches (CountSketch and SJLT) [CW17],
  - Trigonometric sketches (SRHT and SRTT) [AC09],

the best known rate is  $\left(C \log(dT/\delta) \cdot \frac{d}{m}\right)^T$  [PW17].

<u>Note</u>: Extra constant and logarithmic factors in the bound means no analytic expressions for the optimal step size  $\mu_t$ 

### Subspace embedding

(most prior work)

- Standard approximation guarantee for sketching methods
- Leads to subopotimal convergence rates:  $\left(C \log(dT/\delta) \cdot \frac{d}{m}\right)^T$

$$\mathbf{A}_f(\widetilde{\mathbf{x}}_t)^{\top} \mathbf{S}_t^{\top} \mathbf{S}_t \mathbf{A}_f(\widetilde{\mathbf{x}}_t) \approx_{\eta} \nabla^2 f(\widetilde{\mathbf{x}}_t).$$

### Method of inverse moments

(this work)

- Originally proposed for quadratic objectives [LP19]
- Leads to precise convergence rates and optimal step sizes
- Requires inverse moments of the sketched Hessian

$$\mathbb{E}\left[\left(\mathbf{A}_{f}(\widetilde{\mathbf{x}}_{t})^{\top}\mathbf{S}_{t}^{\top}\mathbf{S}_{t}\mathbf{A}_{f}(\widetilde{\mathbf{x}}_{t})\right)^{-k}\right] \quad \text{for} \quad k = 1, 2$$

- Subspace embedding
- 2 Method of inverse moments





### *i*-th leverage score: $\ell_i(\mathbf{A}) = i$ -th diagonal entry of $\mathbf{A}\mathbf{A}^{\dagger}$

- Subspace embedding
- 2 Method of inverse moments



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Uniform Sparsification [CW13]

*i*-th leverage score:  $\ell_i(\mathbf{A}) = i$ -th diagonal entry of  $\mathbf{A}\mathbf{A}^{\dagger}$ 

× x

- Subspace embedding
- 2 Method of inverse moments





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## Implementing LESS Embeddings

- Worst-case implementation (LESS)
  - Preprocessing cost: O(nnz(A) log n + d<sup>3</sup> log d) Approximating leverage scores l<sub>i</sub>(A) [DMIMW12]
  - Sketching cost:  $O(md^2)$ Sparse matrix multiplication **SA**

$$Cost = O(nnz(\mathbf{A})\log n + md^2)$$

### Practical implementation (LESS-uniform)

- Use a uniformly sparsified sketch with  $\alpha d$  non-zeros per row
- If  $\alpha \geq \frac{n}{d} \max_{j} \ell_{j}(\mathbf{A})$ , then we recover theoretical guarantees

$$Cost = O(\alpha m d^2)$$

 $nnz(\mathbf{A}) = number of non-zeros in matrix \mathbf{A}.$ 

### Experiments: Quadratic objective



(b) WESAD dataset

We use sketch size m = 4d, and LESS-uniform has d non-zeros per row.

### Experiments: Logistic regression



We use sketch size m = d/2. Bottom plots report the CPU and GPU wall-clock times to reach a  $10^{-6}$  accurate solution.

- Newton-LESS: Sparsification without trade-offs
  - In Per-iteration efficiency of Sub-Sampled Newton
  - **2** Same convergence rate as Gaussian Newton Sketch
- Sparse sketching can beat Sub-Sampling...
  - In in real-world optimization tasks
  - 2 ... on a variety of hardware platforms
- LESS Embeddings: Fast Gaussian-like sketches
  - Correcting the bias in distributed optimization [DLDM21]
  - 2 Precise convergence rates and optimal step sizes (this work)

Code available at: https://github.com/lessketching/newtonsketch

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